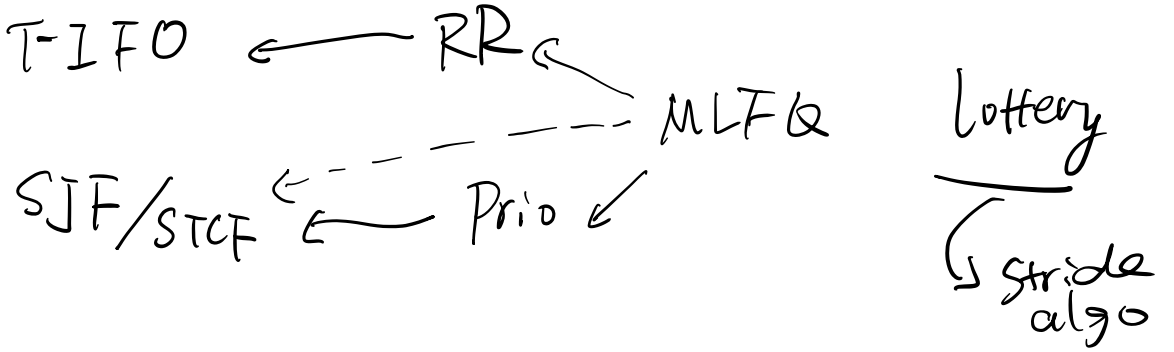
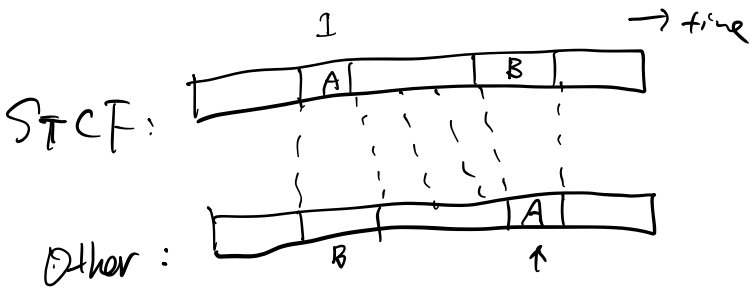
1. Last time
2. More about scheduling
3. Scheduling problem today
4. Lessons and conclusions
5. Threads
6. Intro to concurrency





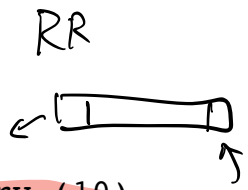
Votes last time:
(with candidate >=5 votes)

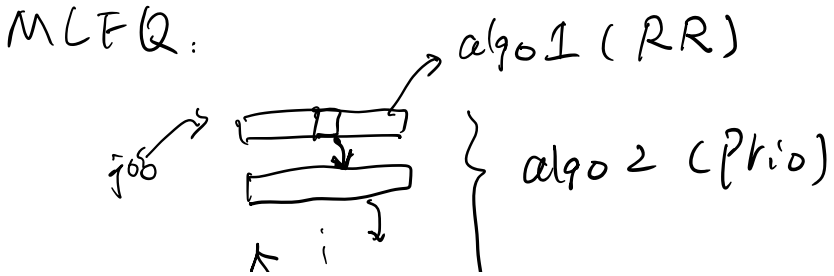
"Best Turnaround Time": STCF (48)

"Best Response Time": RR (34), MLFQ (10)

"Best Fairness": RR (33), MLFQ (10), lottery (10)

"Most popular algorithm": MLFQ (23), lottery (14), STCF (5), Prio (5)





random.

- randomized algo
- SGD
- PoW. PoS

P1 P2 P1 P1 P2

Stride:

P1 (20) cost 1

P2 (10) 2

P1		
0	+1	+1 +1
P2		
0	+2	+2

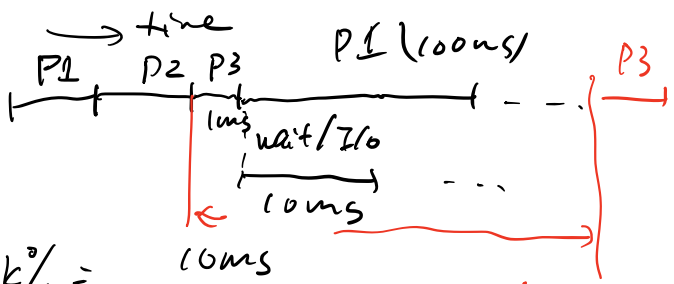
- I/O

3 jobs
  P1, P2: both CPU bound, run for a week
  P3: I/O bound, loop
      (1 ms of CPU, 10 ms of disk I/O)

process	arrival	running
P1	0	1 week
P2	0 +ε	1 week
P3	0 +2ε	30 sec (with 300sec I/O)

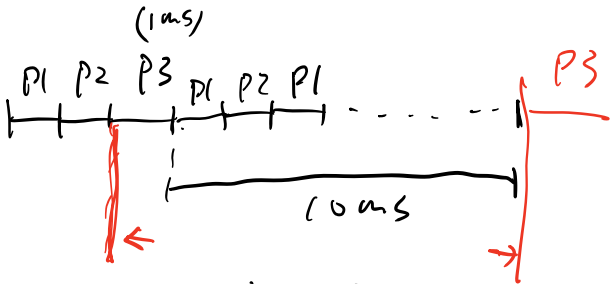
FIFO:

RR (100ms)



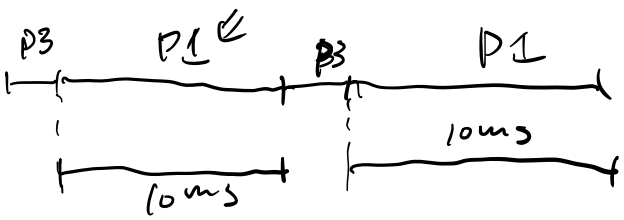
$$\text{disk\%} = \frac{10\text{ms}}{2\times 100\text{ms} + 1\text{ms}} \approx 5\%$$

RR (1ms)



$$\text{disk\%} = \frac{10\text{ms}}{11\text{ms}} \approx 90\%$$

STCF:



$$\text{disk\%} \approx 90\%$$

- predict future EWMA

$t_n \leftarrow n_{th}$ running

$\tau_{n+1} = \alpha \cdot t_n + (1-\alpha)\tau_n$

↳ $\alpha \in (0, 1]$

$\alpha = 1$: $\tau_{n+1} = t_n$

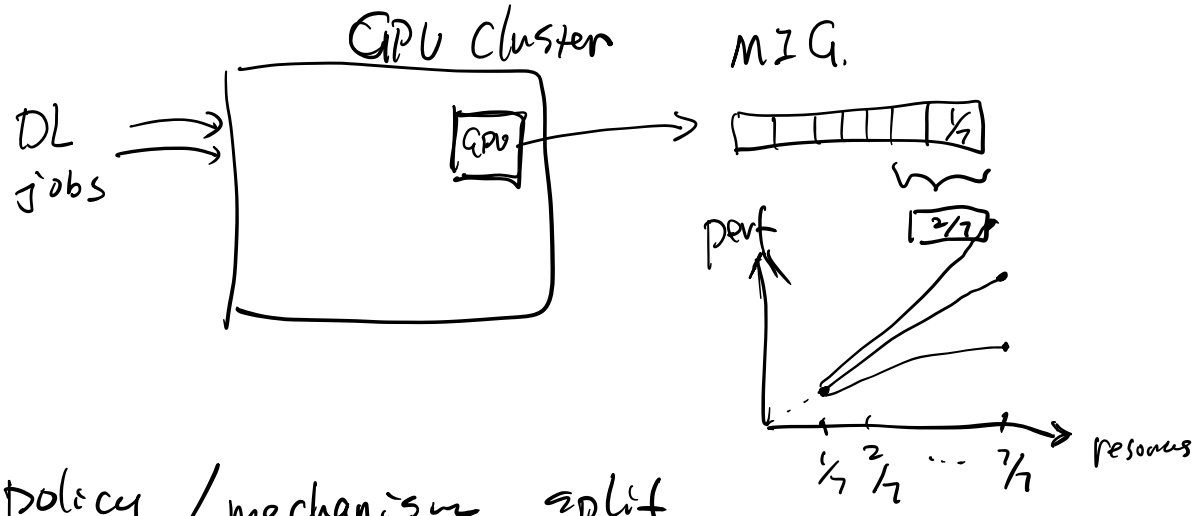
$\alpha = \frac{1}{2}$: $\tau_{n+1} = \frac{1}{2} t_n + \frac{1}{2} \cdot \tau_n$

$(\frac{1}{2} t_{n-1} + \frac{1}{2} \tau_{n-1})$

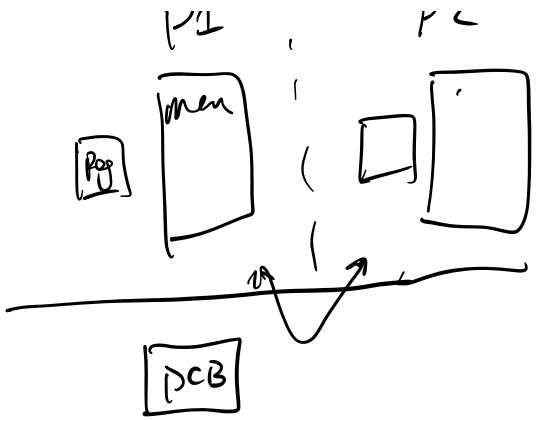
process, context → NN → t

$= \frac{1}{2} t_n + \frac{1}{4} t_{n-1} + \frac{1}{4} \tau_{n-1}$

$+ \frac{1}{8} t_{n-2} + \frac{1}{8} \tau_{n-1}$

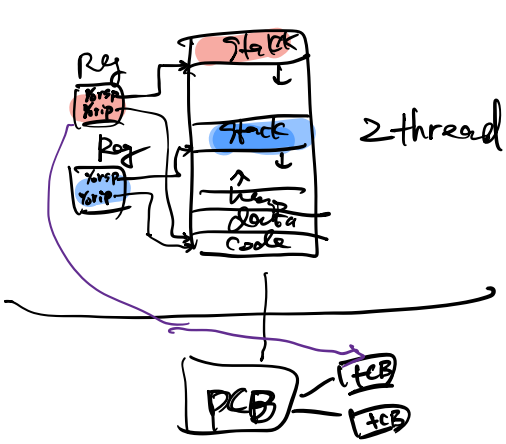


policy / mechanism split

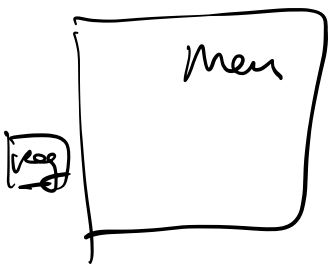


multiple CPU

Thread ⟷ Process



2 threads



thread	process
thread_create(void *fn, ...);	fork() ; ProcessCreate()
thread_exit();	exit()
thread_join(tid);	wait()

A. multiple threads

B. single thread