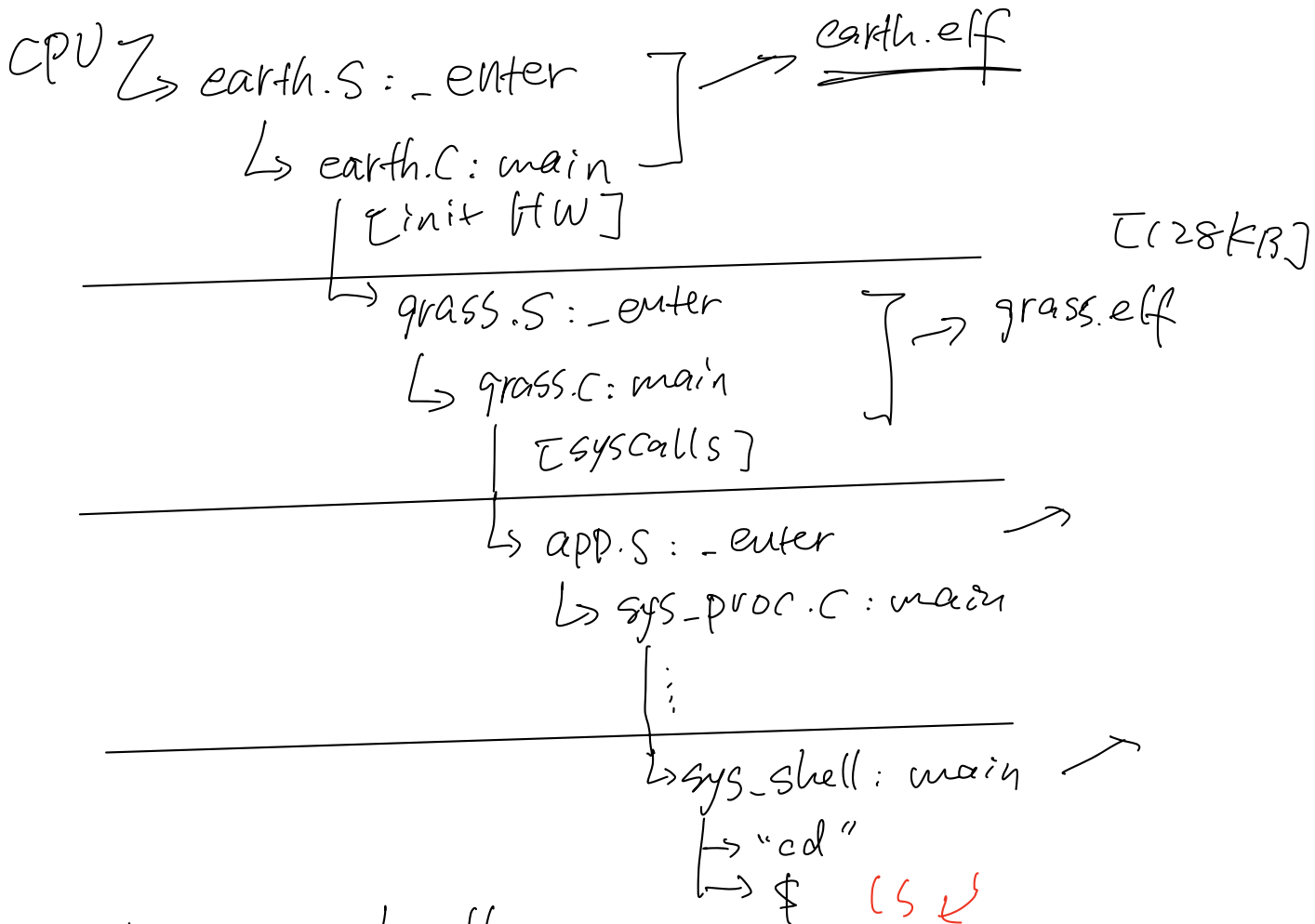


1. kernel \approx three handlers
 2. egos syscall implementation
 3. egos exception handling
-

pid & idx / idx & pid \leftarrow #define . . .

Last time:



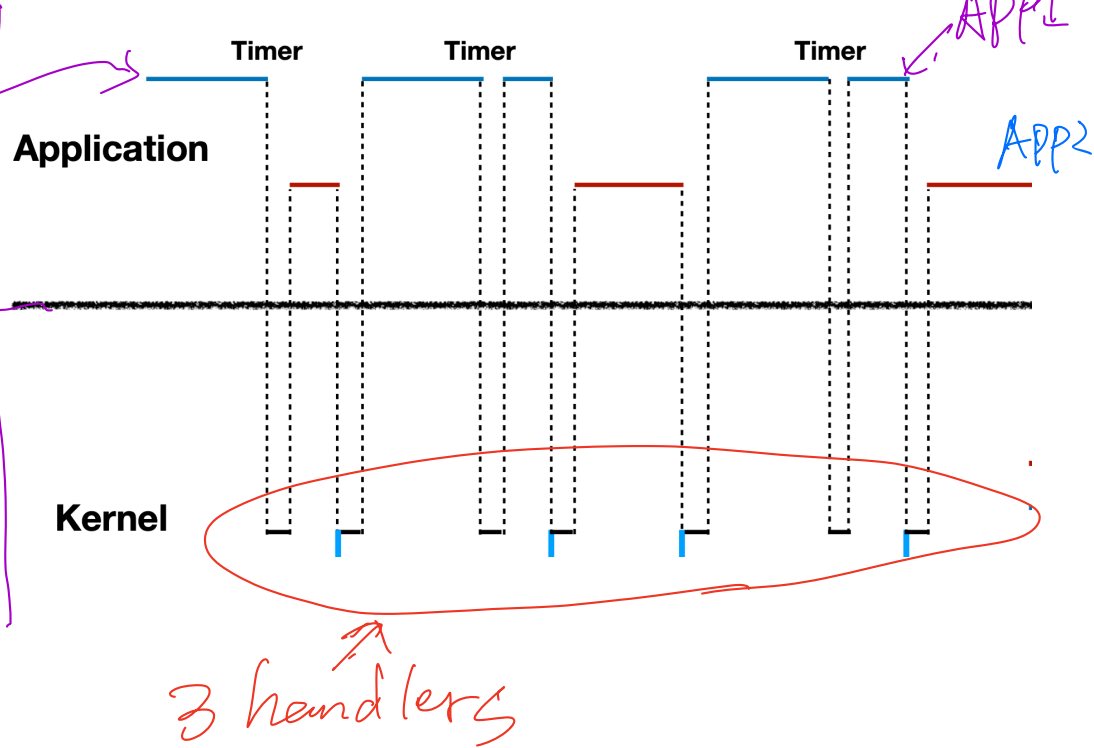
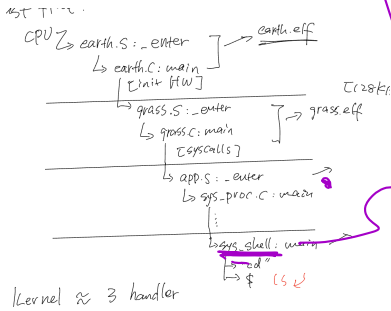
1. kernel \approx 3 handler

Q: What're the 3 ways to trap to kernel?
 interrupts, exceptions, syscalls

a) interrupts.

Executing APPS

Booting



Q: How does a CPU know what to execute when a interrupt is triggered?

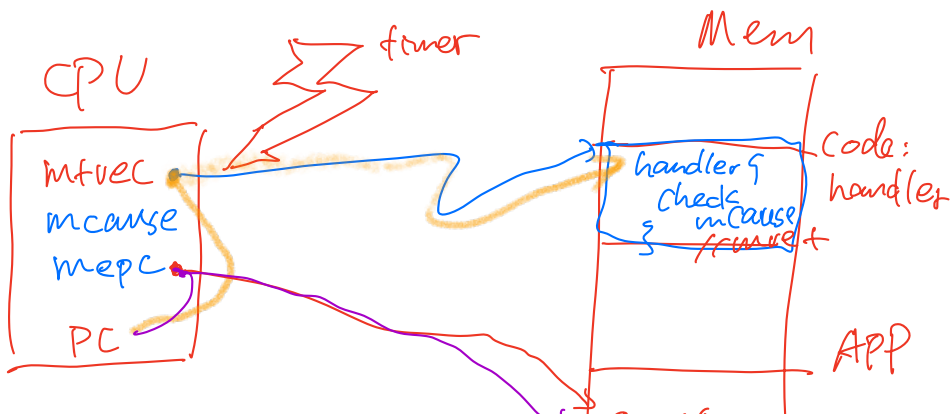
mtvec

Q: How does kernel know which interrupt has been triggered?

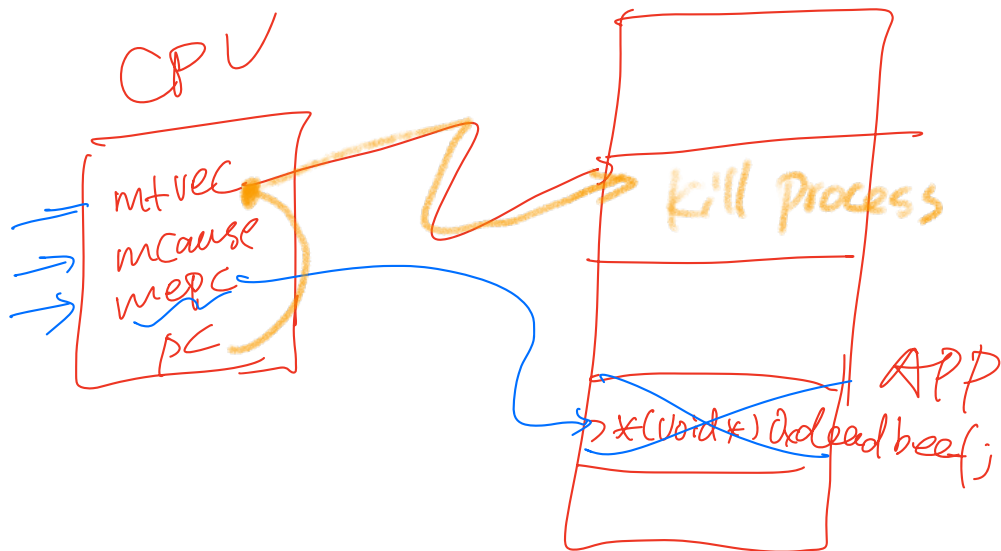
mcause

Q: After handling interrupts, where does CPU ~~to~~ jump to?

mepc

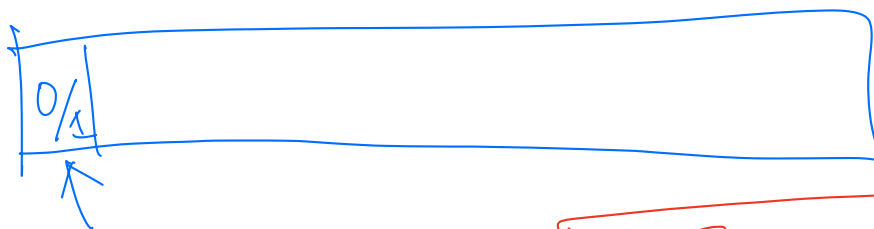


b) exceptions.



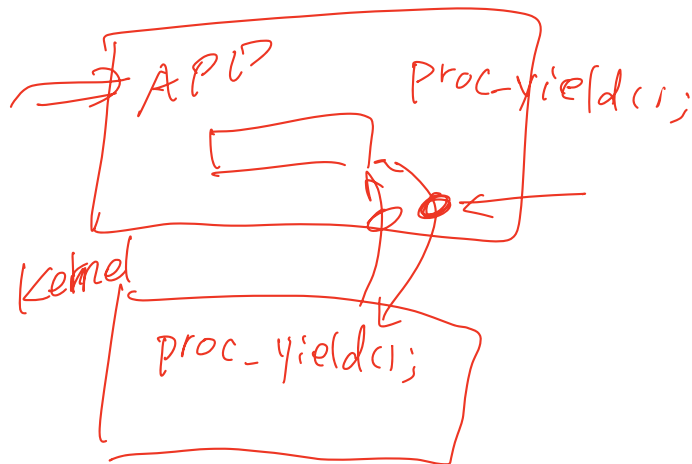
Q: interrupts vs. ~~exe~~ exceptions?

mcause (32bits)



c) syscalls.

x86
int 3



Q1: how to trap to kernel?

- software interrupt
- `*(void*) 0xdeadbeef;`

Q2: how does kernel understand what the application wants?

That is, what information is needed for handling a syscall?

- syscall type
- arguments
- address to return to (mepc)
- return val

Q3: Where to store the information?

- registers
- well-known memory
- stack

2. egos-2k+ syscall implementation

Q1: `sys_inoke()`

Q2: `struct syscall` (^{grass/}`syscall.h`)

Q3: `SYSCALL-ARG`

3. exception handling.

Q: How much Mem "malloc" can get?

- 8GB
- 16GB (2)
- 320GB
- 100GB (2)

170GB

