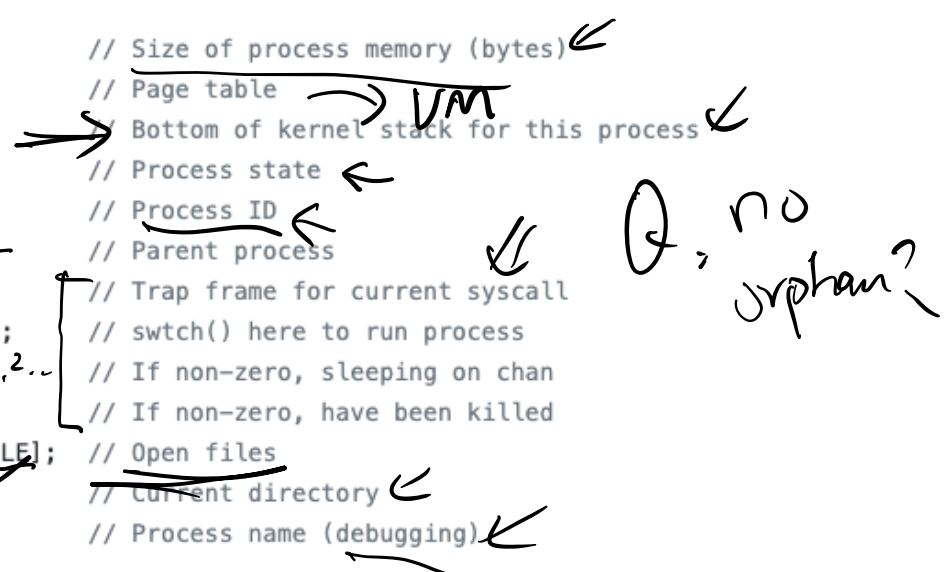


```

36
37 // Per-process state
38 struct proc {
39     uint sz; // Size of process memory (bytes) ←
40     pde_t* pgdir; // Page table ←
41     char *kstack; // Bottom of kernel stack for this process ←
42     enum procstate state; // Process state ←
43     int pid; // Process ID ←
44     struct proc *parent; // Parent process ←
45     struct trapframe *tf; // Trap frame for current syscall ←
46     struct context *context; // switch() here to run process ←
47     void *chan; // If non-zero, sleeping on chan ←
48     int killed; // If non-zero, have been killed ←
49     struct file *ofile[NOFILE]; // Open files ←
50     struct inode *cwd; // Current directory ←
51     char name[16]; // Process name (debugging) ←
52 };
53

```



Borrowed from xv6 <https://github.com/mit-pdos/xv6-public/blob/eeb7b415dbcb12cc362d0783e41c3d1f44066b17/proc.h>

Week 4.a
CS5600
1/30 2023
<https://naizhengtan.github.io/23spring/>

1. Shell internal continued & discussions
 2. Implementation of processes
 3. Context switch intro
 4. Scheduling intro
-

02/05, 02/17

Q: fork() ret val? ← child ret = 0

parent ret < pid of child

Q: fd 0/1/2? ← 0: stdin
① stdout

Q: "ls > log.txt"? 2: stderr

1 CS5600 23spring
 2 Handout week03b
 3
 4 The handout is meant to:
 5
 6 --illustrate how the shell itself uses syscalls
 7
 8 --communicate the power of the fork()/exec() separation
 9
 10 --give an example of how small, modular pieces (file descriptors,
 11 pipes, fork(), exec()) can be combined to achieve complex behavior
 12 far beyond what any single application designer could or would have
 13 specified at design time.

1. Pseudocode for a very simple shell

```

17 while (1) {
18   write(1, "$ ", 2);
19   readcommand(command, args); // parse input
20   if ((pid = fork()) == 0) { // child?
21     execve(command, args, 0);
22   } else if (pid > 0) { // parent?
23     wait(0); //wait for child
24   } else {
25     perror("failed to fork");
26   }
27 }

```

2. Now add two features to this simple shell: output redirection and backgrounding

By output redirection, we mean, for example:
 \$ ls > list.txt
 By backgrounding, we mean, for example:
 \$ myprog &
 \$

```

37 while (1) {
38   write(1, "$ ", 2);
39   readcommand(command, args); // parse input
40   if ((pid = fork()) == 0) { // child?
41     if (output_redirected) {
42       close(1);
43       open(redirect_file, O_CREAT | O_TRUNC | O_WRONLY, 0666);
44     }
45     // when command runs, fd 1 will refer to the redirected file
46     execve(command, args, 0);
47   } else if (pid > 0) { // parent?
48     if (foreground_process) {
49       wait(0); //wait for child
50     }
51   } else {
52     perror("failed to fork");
53   }
54 }
55

```

3. Another syscall example: pipe()

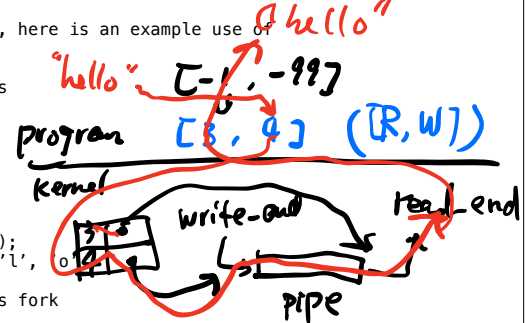
58 The pipe() syscall is used by the shell to implement pipelines, such as
 59 \$ ls | sort | head -4
 60 will see this in a moment; for now, here is an example use of pipes.

63 // C fragment with simple use of pipes

```

64
65 int fdarray[2]; [-1, -99]
66 char buf[512];
67 int n;
68
69 pipe(fdarray);
70 write(fdarray[1], "hello", 5);
71 n = read(fdarray[0], buf, sizeof(buf));
72 // buf[] now contains 'h', 'e', 'l', 'l', 'o'
73

```



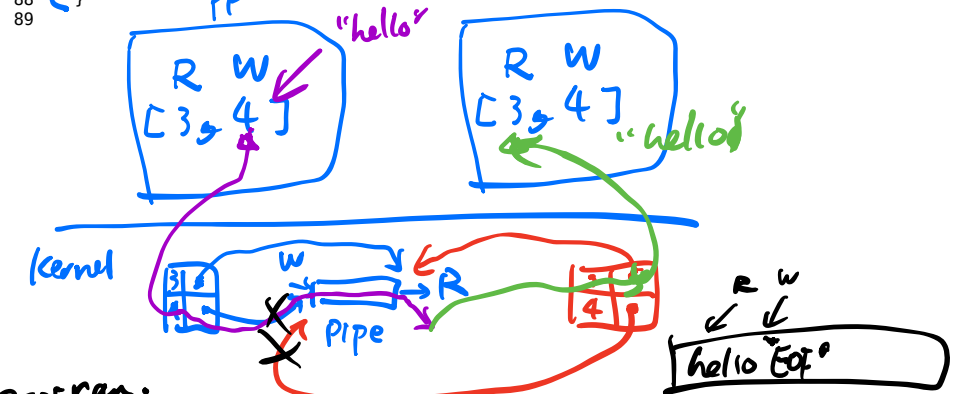
4. File descriptors are inherited across fork

76 // C fragment showing how two processes can communicate over a pipe

```

77
78 int fdarray[2];
79 char buf[512];
80 int n, pid;
81
82 pipe(fdarray);
83 pid = fork();
84 if (pid > 0) { // parent
85   write(fdarray[1], "hello", 5);
86 } else { // child
87   n = read(fdarray[0], buf, sizeof(buf));
88 }
89

```



Program:
 while (1) read (fd, ...) &

EOF

close(pipefd[0]);

• fork / exec. Separation

```

if( fork() == 0 ) { //child
    // "ls"
    exec( binary, ... ) ←
} else { //parent
    ...
}

```

ϕ mem \rightarrow fork
 Redirection / Pipe / ...
 PM arg = ...



Environment

--syscall CreateProcess on Windows:

```

BOOL CreateProcess(
    name,
    commandline, ← "ls"
    security_attr, ← "a"
    thr_security_attr,
    inheritance?,
    other flags,
    new_env,
    curr_dir_name,
    ...)

```

• good abstraction ?

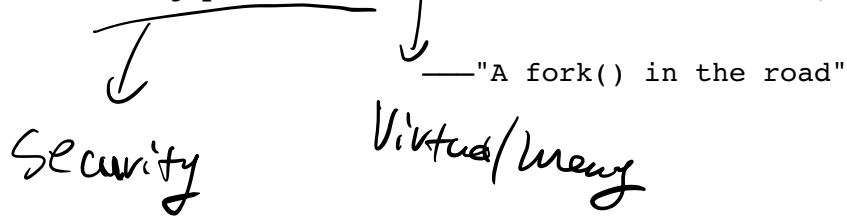
- fd \leftarrow open/close/read/write
- fd
- o/p/z
- fork/exec.

• "A fork() in the road"

HotOS'19

↑ on current world

"Fork today is a convenient API for a single-threaded process with a small memory footprint and simple memory layout that requires fine-grained control over the execution environment of its children but does not need to be strongly isolated from them. In other words, a shell."



"Fork doesn't compose" ←

— example:

```

print("hello world");
fork();
print("\n");
    
```

2 processes

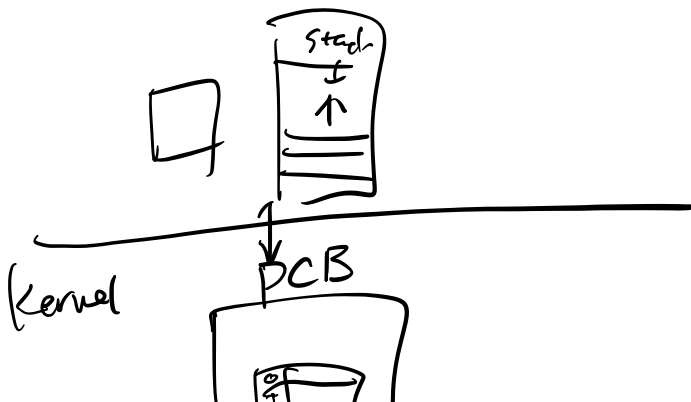
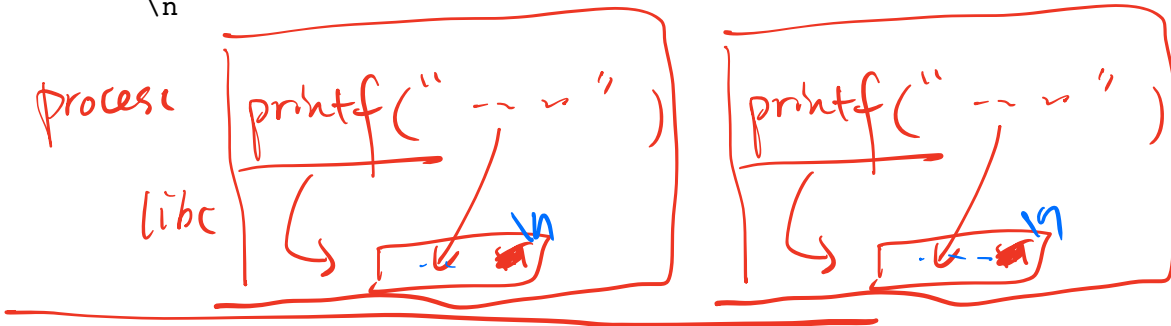
hello... \n
 \n ←
 \n ←

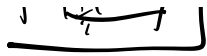
QUESTION: what do you expect to see on screen?

A:
 hello world\n
 hello world\n

B:
 hello world\n
 \n

3
 A or B
 |||





- Context Switch

