

CS 3650 – Computer Systems
Spring 2024
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Lecture 15, Tue Feb 27 2024

Midterm exam logistics

conflicts: wed 2:50-4:30 Cavall 097
proctor will have exam w/ your name
→ no notes or other material

section 3 exam - regular time / place
your own handwritten notes
(iPad etc? print them)

Exam review

- * C programming language
 - variable scope -
 - arguments passed by value

int a = 5;

f(a) ← copies 'a'
(passes 5)

g(x) — — — — — f(int x)

function can take
address arg and
change things that arg points to

g(int *x) {
*x += 1;
}

argument:

f(int a) {
↑
'a' range of validity
}

local:

f(..) {
int a;
int b;
↑
'b' range of validity
}

if () {
int x
↑
}

* shell - simple commands & redirection

cd, pwd, echo, cat, ls, mkdir, grep (patterns)

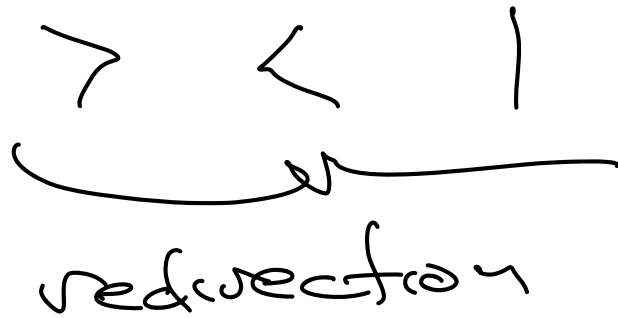
read "input" - stdin
or named files

"find" or "search"

same input style
as cat - e.g.

grep pattern file

ls | grep pattern



wildcards: *

<expression> *
0, 1, ... of <expression>

ls *.c abc*

expands to all files
matches <zero-or-more> of
anything

* file descriptors, dup2, etc

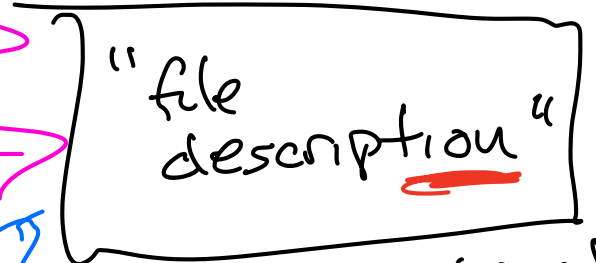
open(name, O_RDONLY)

(cascade kernel)

3 ←
↑
file descriptor

dup2(3, 4)

4



current position

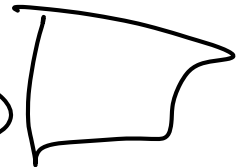
reference count

fork()

child:
3, 4
ref → 4

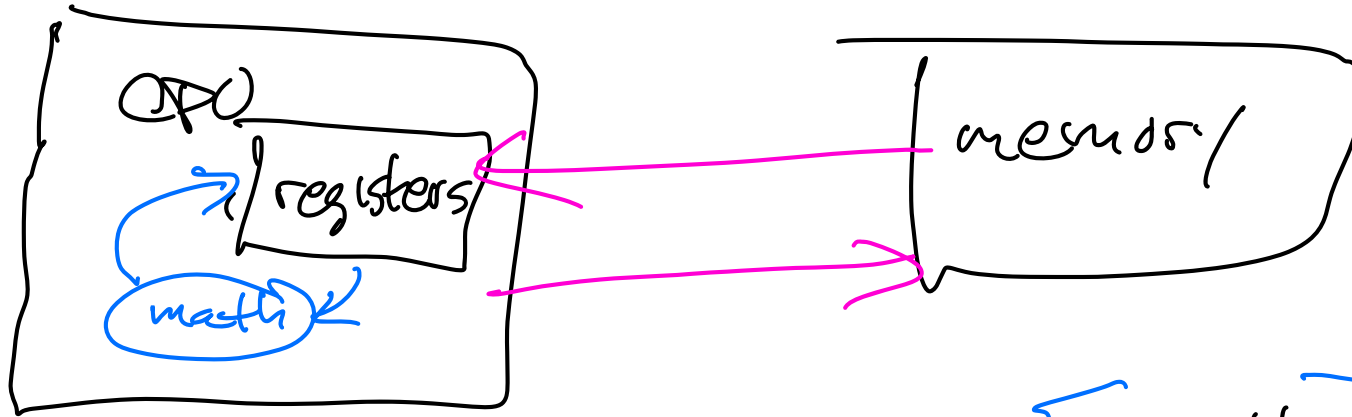
0 →

1 →



2 ways to copy file descriptor (i.e. add reference)
① Fork ② dup

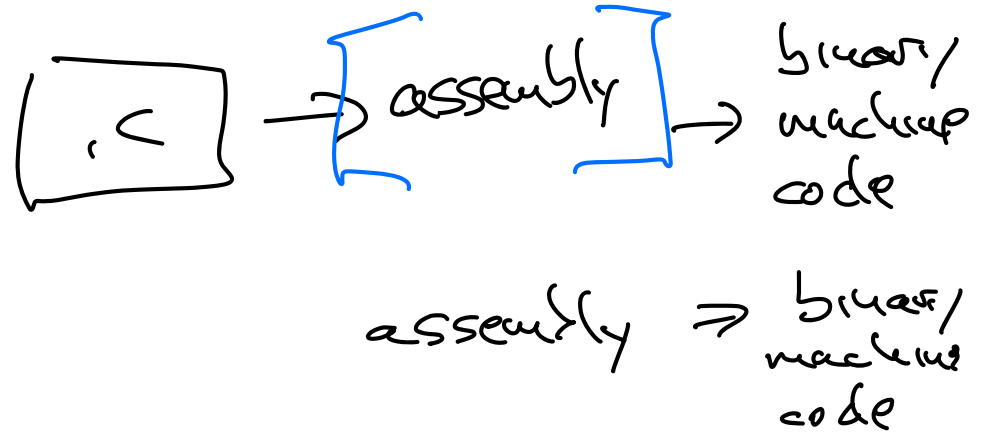
* CPU & registers



read mem → res
res → reg operations
write res → mem

CALL <addr>

RET



* race conditions

deposit (sum) Σ

1) tmp = 0 int tmp = balance

3) tmp = 10 tmp += sum

5) bal = 10 balance = tmp

A

~~30~~

B

30

dep (~~30~~)

0

2) int tmp = balance

tmp = 0

4) tmp = ~~10~~ 30 tmp += sum

balance = tmp

6) bal = ~~10~~ 30

\Rightarrow reason about 2 concurrent executions of same code

* monitors
(mutex + CV)

"Condition" =
queue not
empty

atomic

mutex_lock(u)

(mutex)

while (condition not true)

wait(CV, u)

locks mutex
before returning

queue
not
empty

do something

mutex_unlock(u)

mutex_lock(u)
make condition
true
signal(CV)
unlock(u)

avoiding TOCTTOU

time of check to time of use

wait - monitor

while ~~if~~ condition not true {

unlock

wait

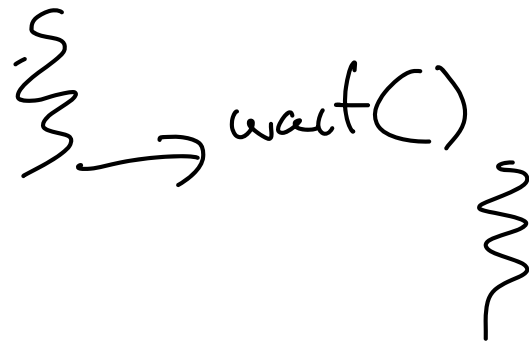
lock

unlock

time of check

make condition T
signal

time of
use



"give the sequence of ..."

deposit (sum) Σ

① int tmp = balance

② tmp += sum

③ balance = tmp

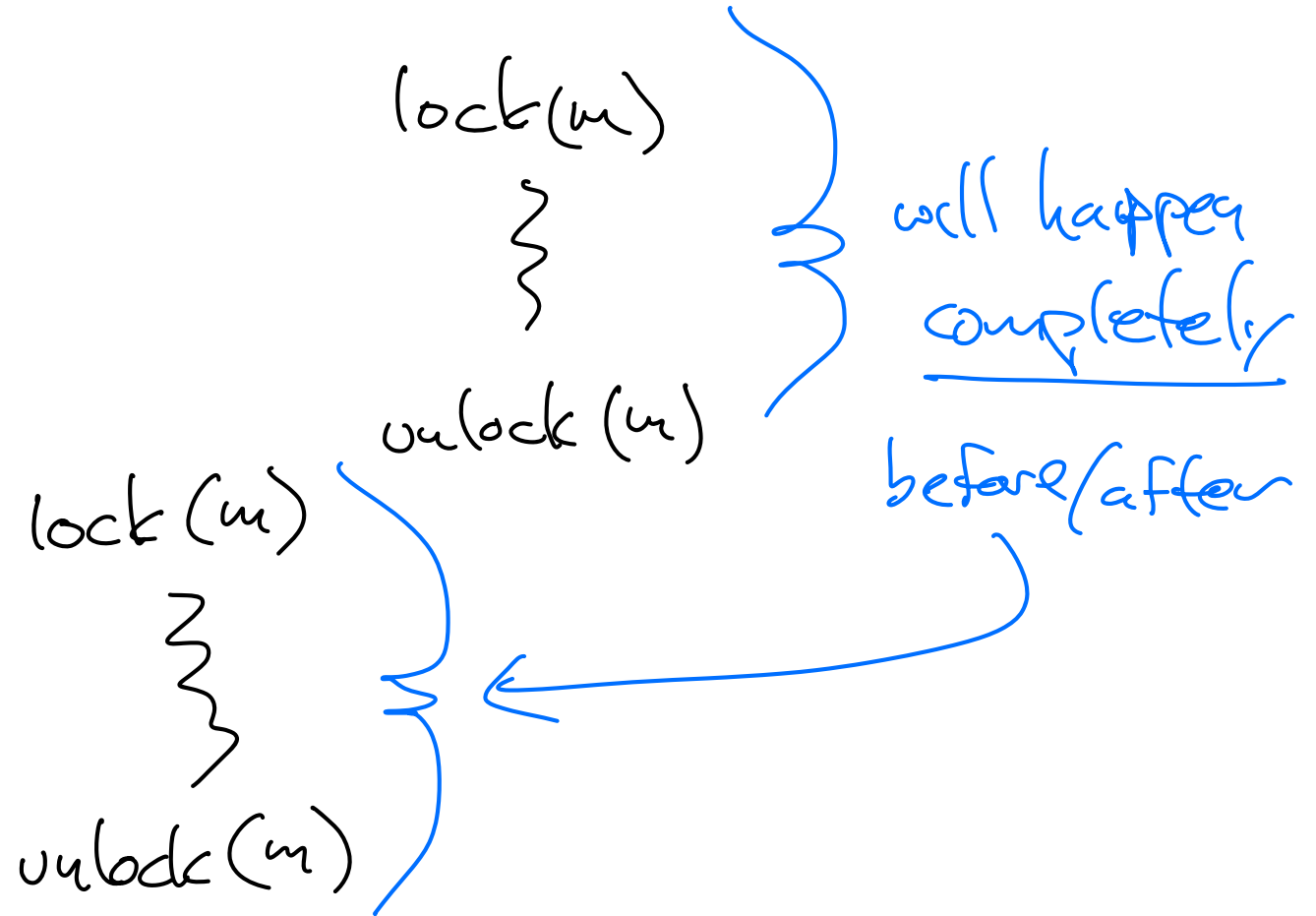
threads A, B

A1 ← thread A executes line 1

B2 ← B 2

A1 A2, A3, B1, B3, B2

6 questions 16 points each + 4 Free points



It's not a monitor if
there's no CV

lock(m)
while !P
signal C
unlock(m)

lock(m)
while !P
wait(C, m) ----- sleep
----- wake up <
-----<