

CS 3650 – Computer Systems
Spring 2024
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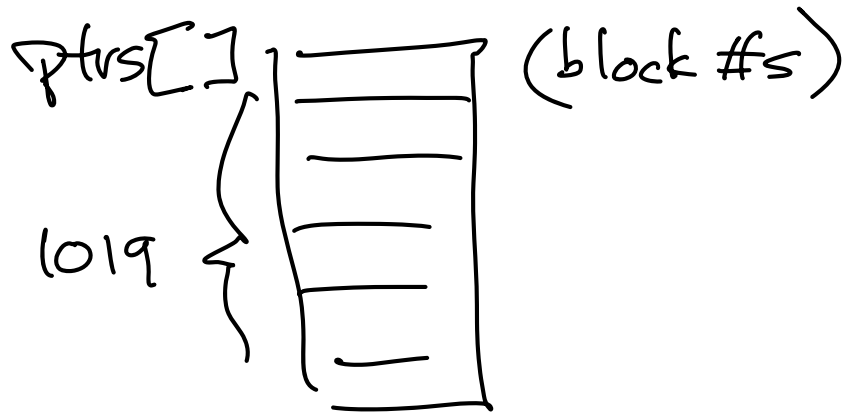
Lecture 20, Thur Mar 21, 2024

Lab 4

4096 bytes

inode:

size (bytes)



1019 · 4096 B

superblock:

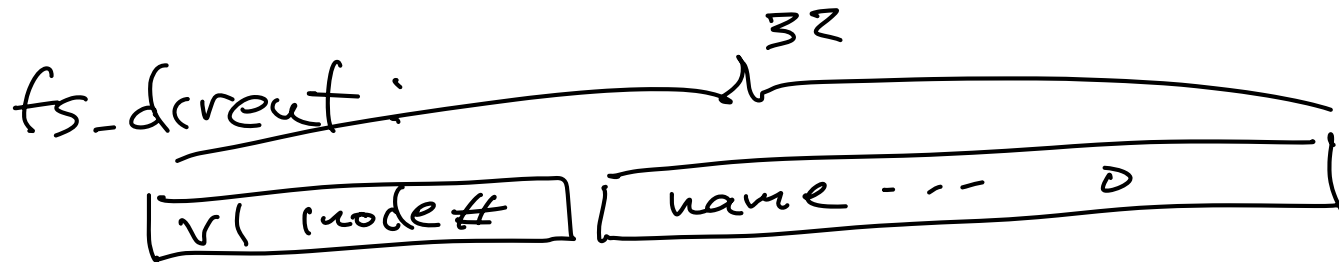
disk size: <n>

root node #: <n>

~~allocation bitmap~~

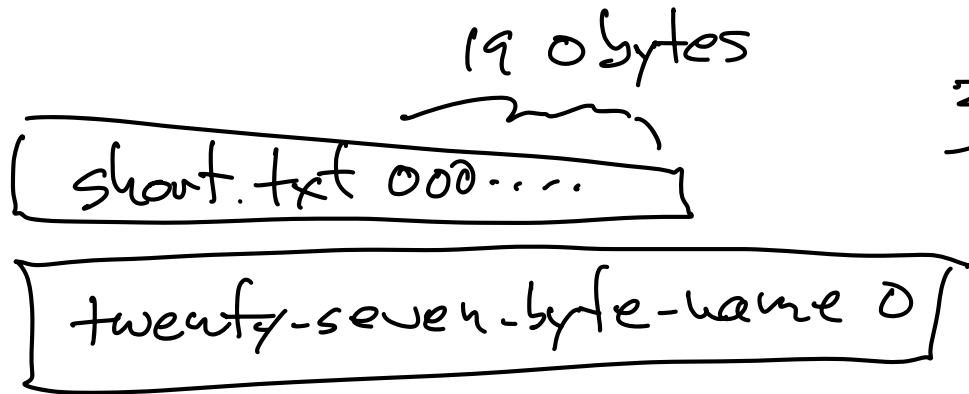
```
struct fs_inode in;  
block_read (&in,  
            block #,  
            1)
```

in.ptrs[0] ←



valid

struct {
 valid
 mode
 name [28]
 }



struct fs_dirent d [28]

block_read (&d, block#, 1)

if (d[0].valid) printf("0x%05: %0d",
 d[0].name, d[0].mode)

FUSE readdir:

```
readdir (char *path,  
         void *ptr,
```

```
int (*filler) (void *ptr,  
              char *name,  
              struct stat *st,  
              ...)
```

for each name F in dir:

~~get attr (F) → sb~~

filler(ptr, F, &sb)

find file attributes

put them in sb

↳ struct stat

```
struct fuse_operations {  
    int (*read) (char *path,  
               char *buf,  
               int len,  
               int offset)  
    int (*getattr) (...)  
    :  
};
```

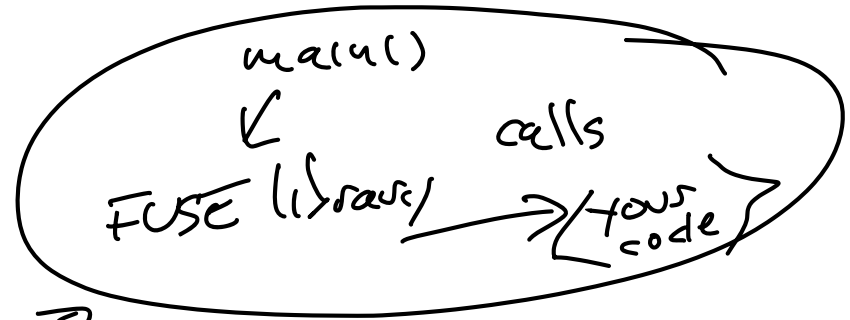
fs_ops.read (.....)

maybe pass
NULL for
stat buf?
check...

How FUSE works

bash / other process

open(), read(), etc



kernel

FUSE

ext4

~
~

getattr(char *path,
struct stat *sb)

factor: label node \rightarrow struct stat

read(char *path,
char *buf, int len,
int offset, ...)

Path translation

int pathlen

char * names[]

→ "home"

→ "pjd"

→ "file.txt"

path → inode #

inode = root_inode

→ for name in names

read inode

check ISDIR

read DE block

if DE[i].valid & name match

inode = DE[i].inode

→ next name

/home/pjd/file.txt

inode #

inode = 1

→ _in = readblk

check ISDIR

read block

, name = "home"

inode = ... inode

DO NOT USE strcpy

memcpy ← copy bytes

(void *dst, void *src, int #bytes)

forget about EOF

syscalls: fd = open
val = read(fd, buf, len)

if val == 0

→ we're at end
of file

read(offset, len)

→ offset > file len: return 0
bytes

stdio: FILE*

int c = getc(fp)

if (c == EOF)

0...255 ←

or -1

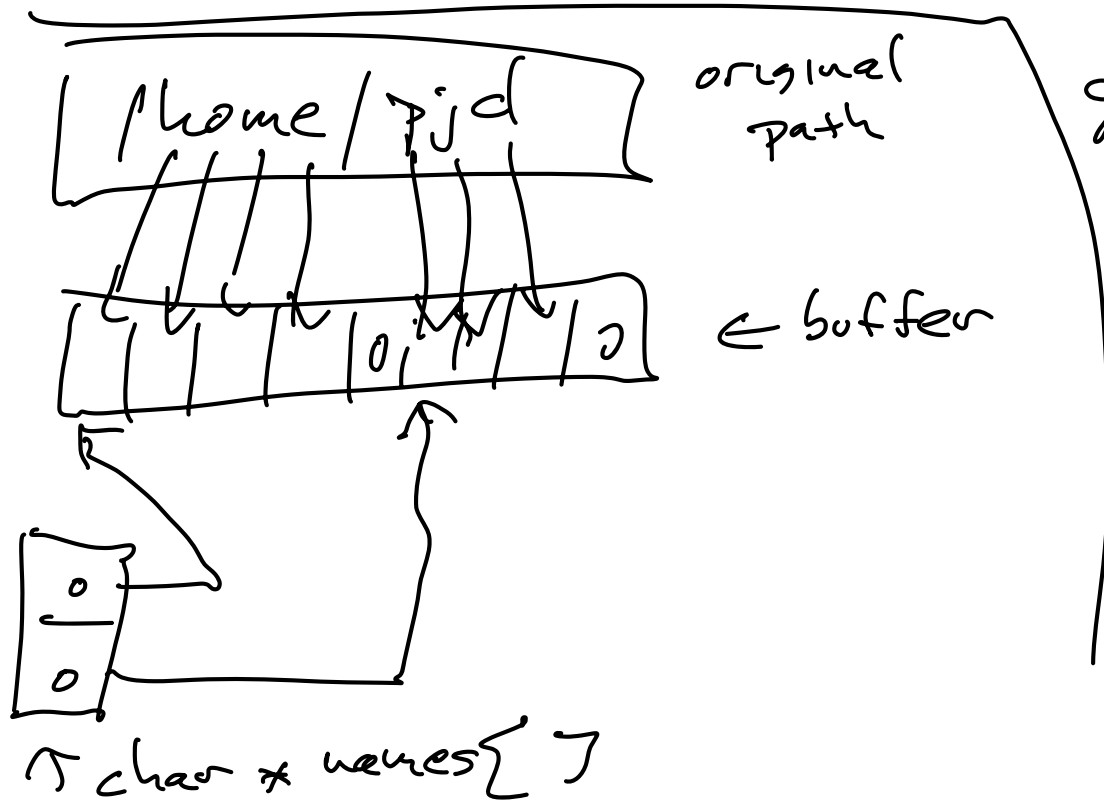
return values :

< 0 = error

≥ 0 = success

read : ≥ 0 : bytes read

if (not found)
return -ENOENT;



getatfr ()

split

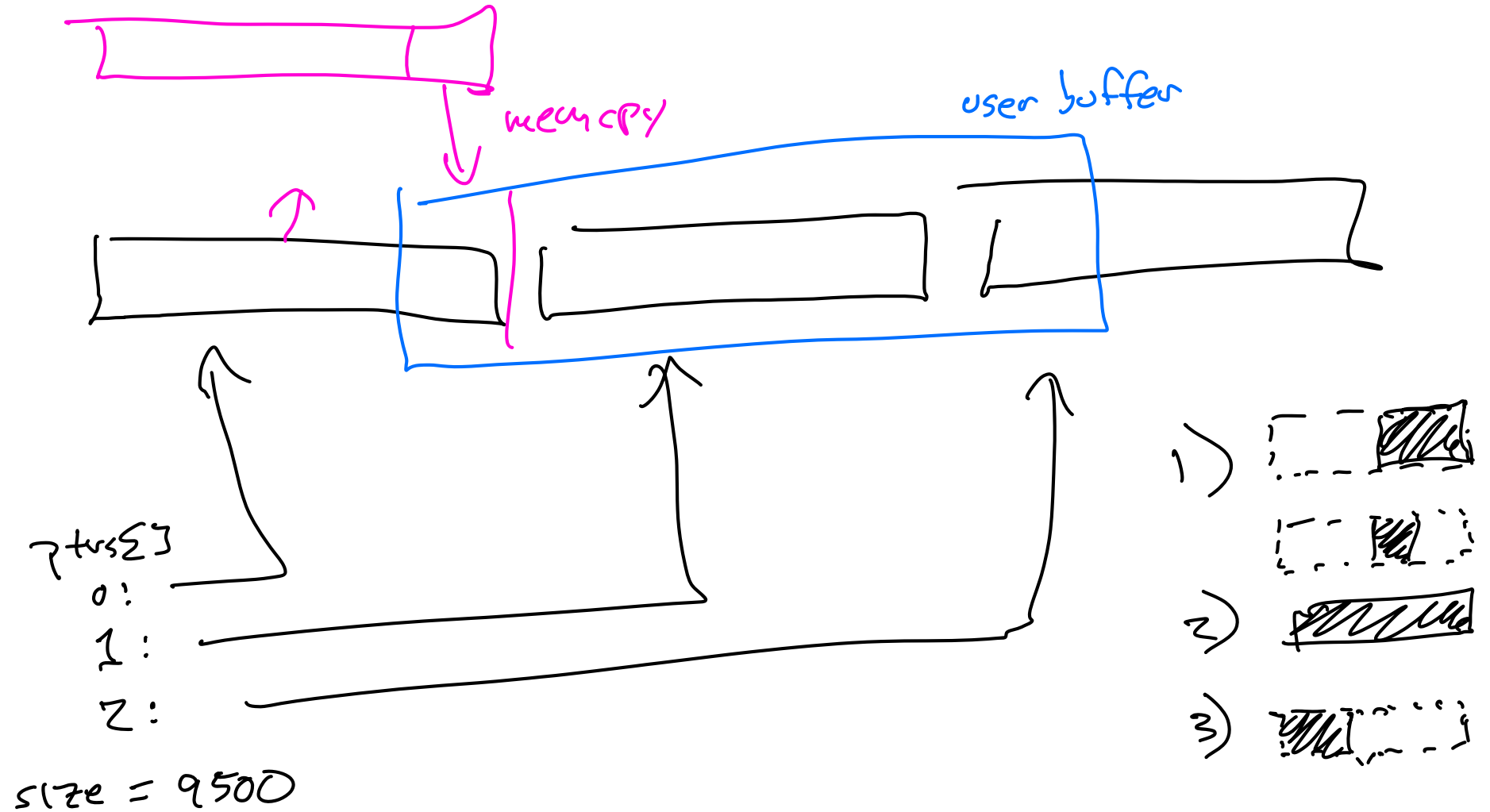
ret i = translate (...)

if i < 0

return i

read

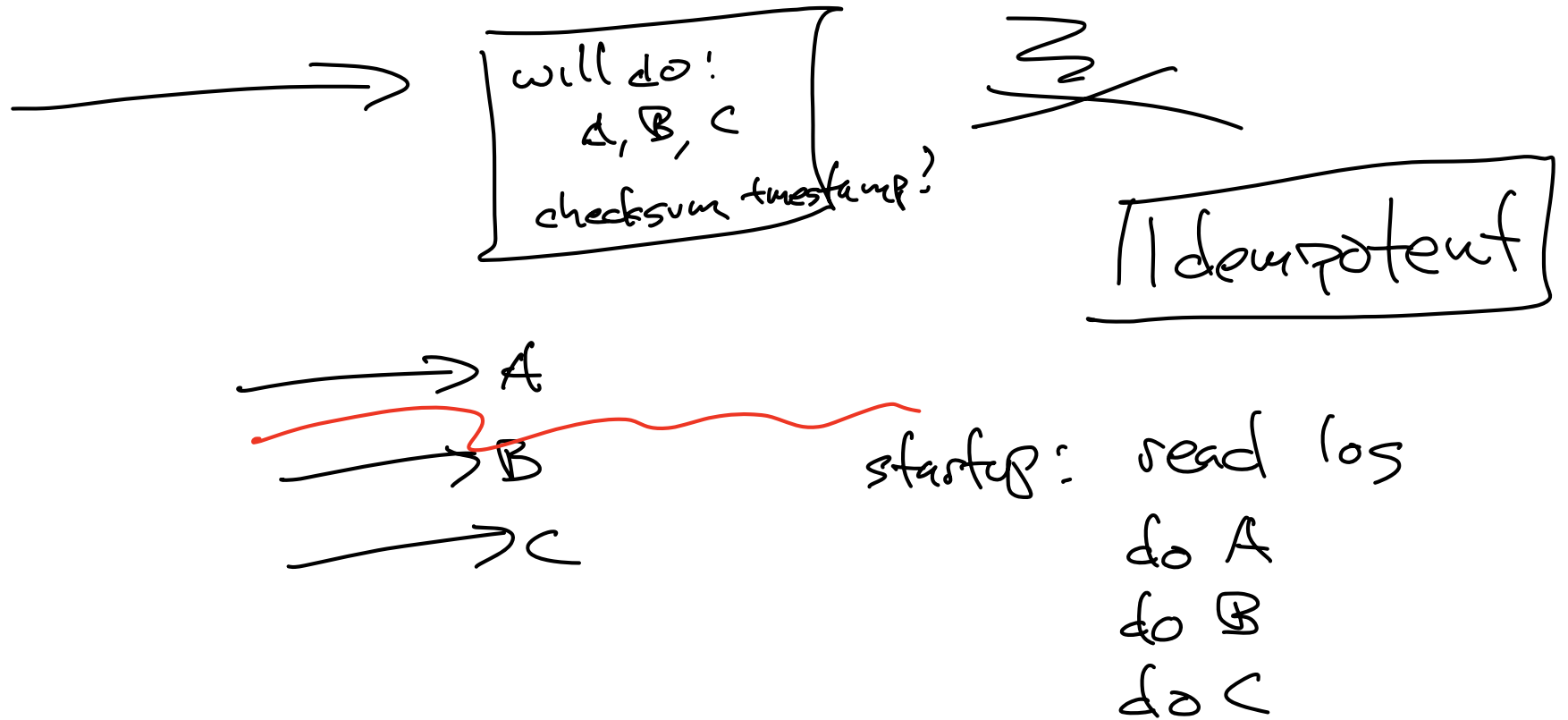
read (offset = 3000, len = 6000)



More file & related stuff

journalling (write-ahead log)
→ crash consistency

similar (sort of)
to multi-thread

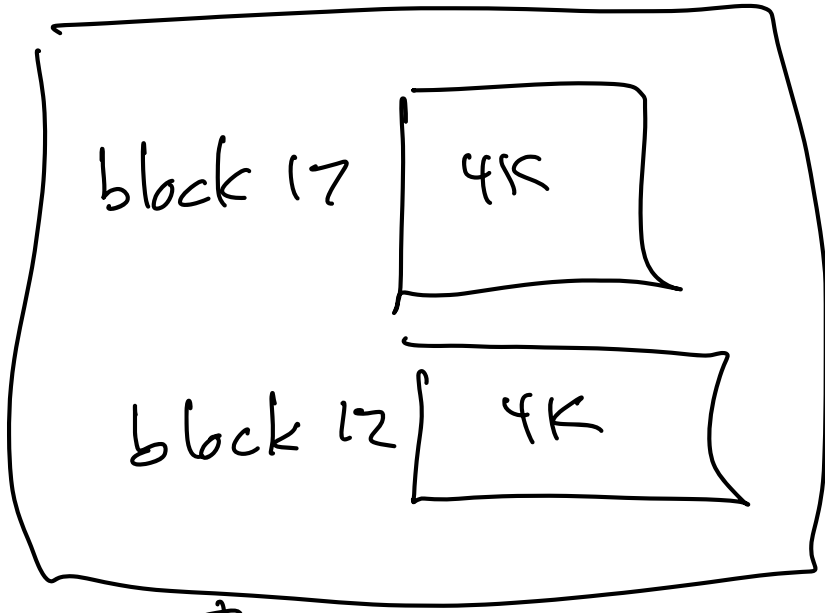


Idempotent operation:

1 time }
N > 1 times } same

$x = 1$ ✓

$x = x + 1$ ✗



↑
journal record

Network File System (NFS)



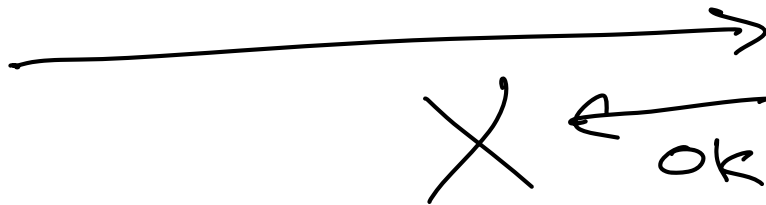
NFS network failure handling

idempotent

READ (fileID, offset, len) →

WRITE (fileID, offset, len, data) →

creat ^{write} → X server



Not file systems

S3

can't:

- modify object
- rename directories

"/dir/dir2/file"

"~~dir~~/dir2/file"

rename dir → Xdir

PUT ^{key} objname
(data)

GET objname

DELETE

+ byte range request

← key