

```
36
37 // Per-process state
38 struct proc {
39     uint sz;           // Size of process memory (bytes)
40     pde_t* pgdir;    // Page table
41     char *kstack;    // Bottom of kernel stack for this process
42     enum procstate state; // Process state ★
43     int pid;          // Process ID
44     struct proc *parent; ● // Parent process
45     struct trapframe *tf; // Trap frame for current syscall
46     struct context *context; █ // swtch() here to run process
47     void *chan;        // If non-zero, sleeping on chan
48     int killed;        // If non-zero, have been killed
49     struct file *ofile[NFILE]; // Open files █
50     struct inode *cwd; // Current directory
51     char name[16];    // Process name (debugging)
52 };
53
```

Borrowed from xv6: <https://github.com/mit-pdos/xv6-public/blob/eeb7b415dbcb12cc362d0783e41c3d1f44066b17/proc.h>

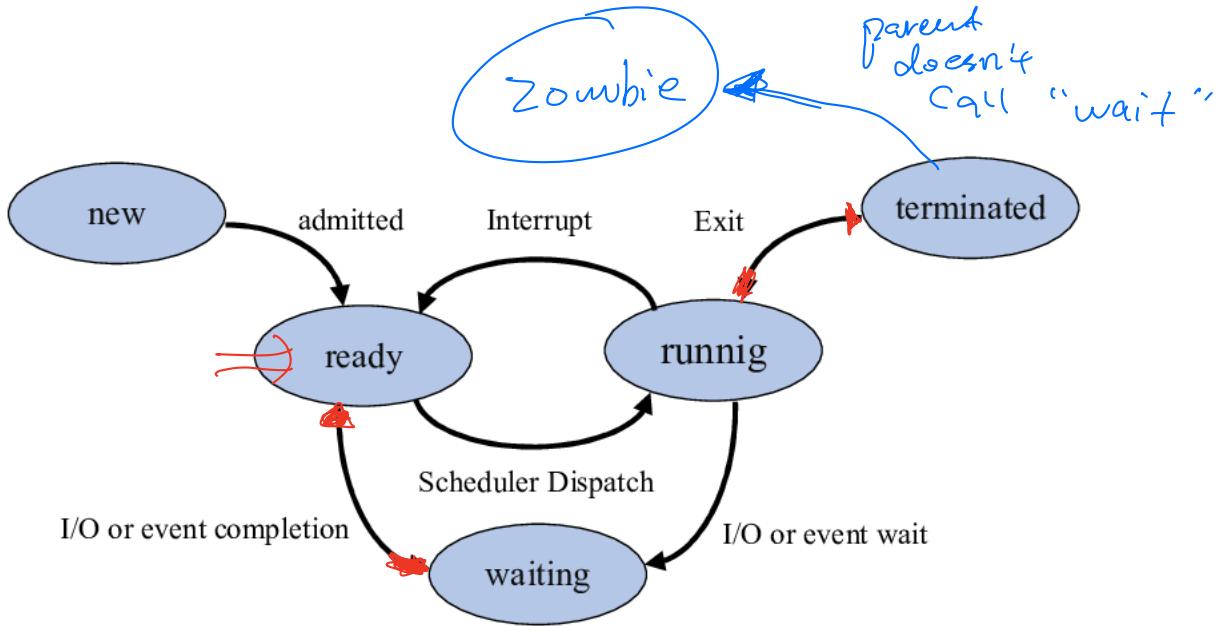
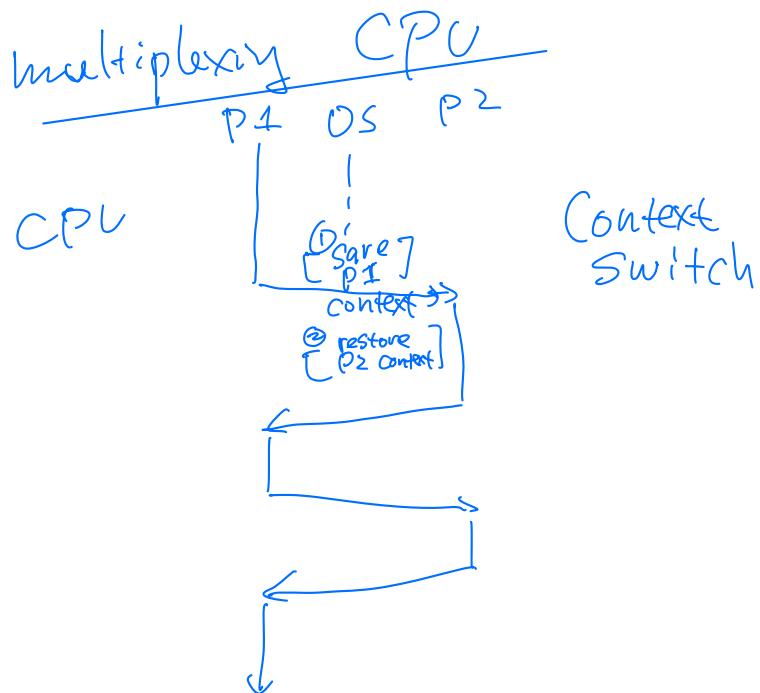
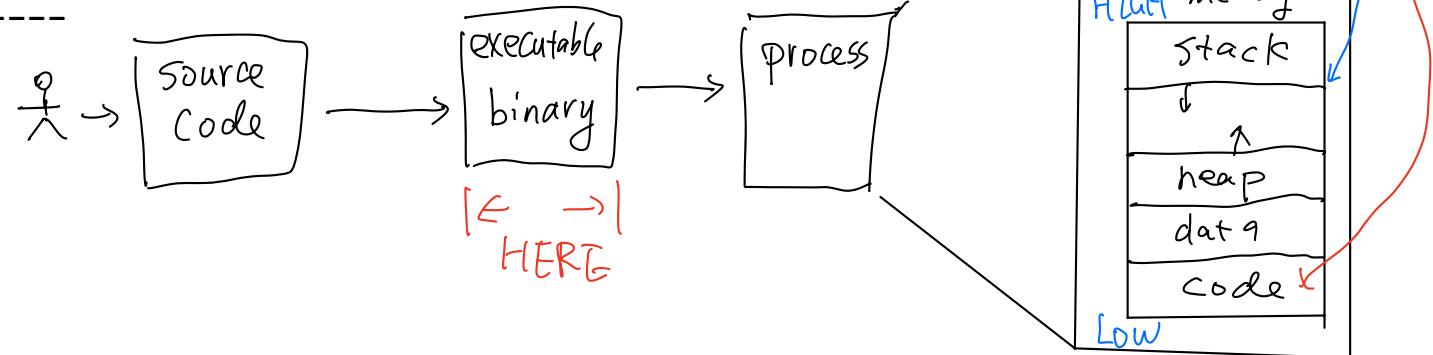


Fig13 from "Priority based round robin (PBRR) CPU scheduling algorithm"

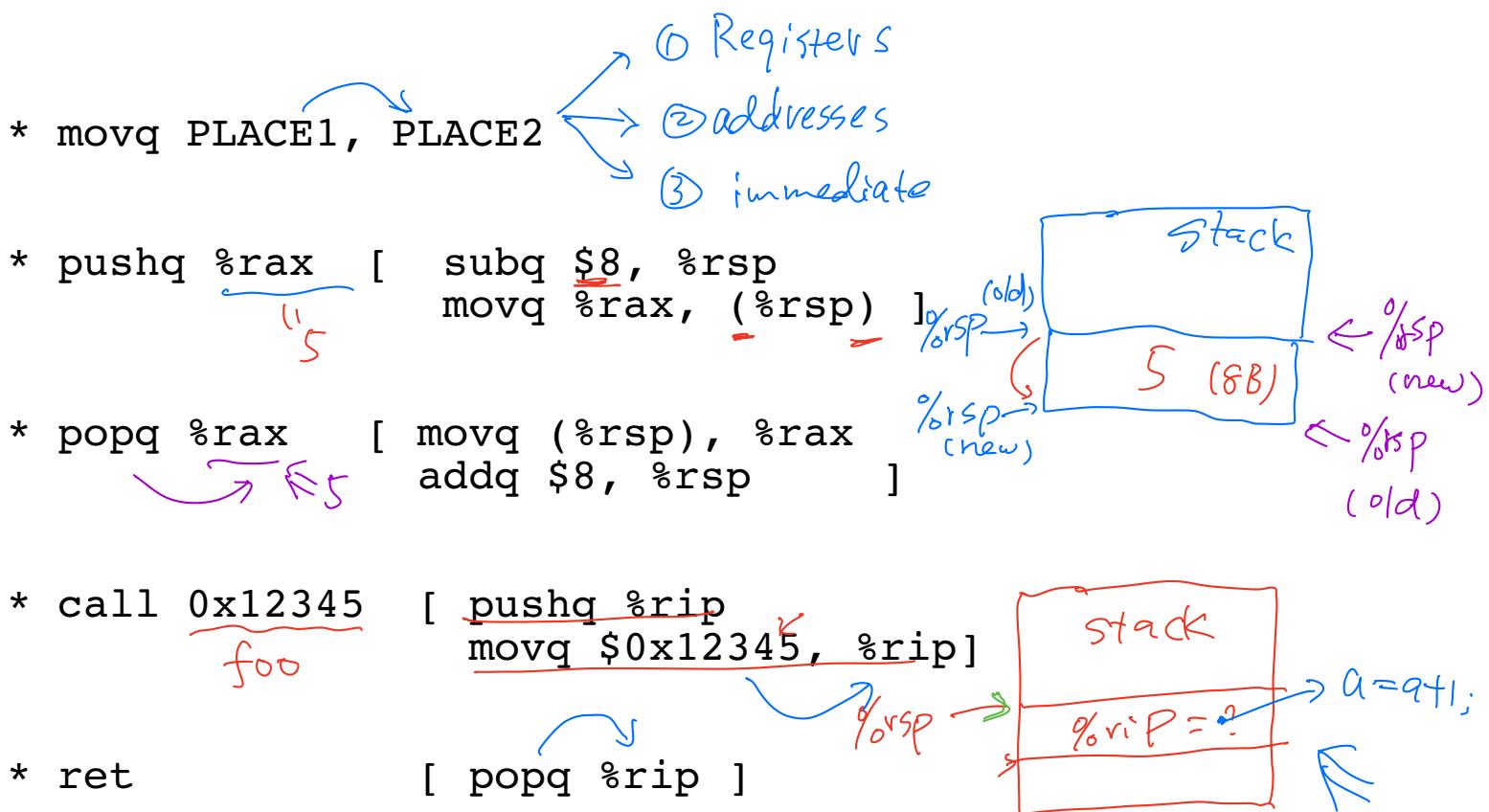


1. x86-64 assembly (cont'd)
2. Stack frames
3. Implementation of processes
4. Context switch intro
5. OS scheduling



Crash course of x86-64 assembly

`int X=5;`



```
main() {
    foo();
    a=at();
}
```

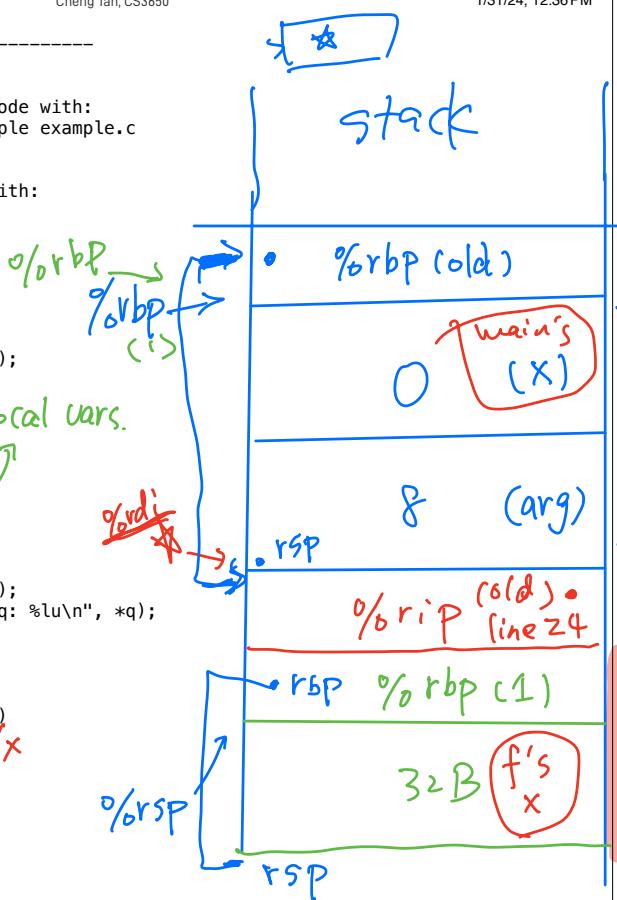
Annotations: %rip is crossed out; a=at(); is highlighted in green; foo(); is highlighted in purple; return; is highlighted in blue; and ?? is highlighted in red.

- %rip : instruction ptr (next insy)
- %rsp : stack ptr
- %rbp : base ptr

-----[example.c]-----

```

1  /* CS3650 -- handout w04b
2  * compile and run this code with:
3  * $ gcc -g -Wall -o example example.c
4  * $ ./example
5  *
6  * examine its assembly with:
7  * $ gcc -O0 -S example.c
8  * $ [editor] example.s
9  */
10
11 #include <stdio.h>
12 #include <stdint.h>
13
14 uint64_t f(uint64_t* ptr);
15 uint64_t g(uint64_t a);
16 uint64_t* q;
17
18 int main(void)
19 {
20     uint64_t x = 0;
21     uint64_t arg = 8;
22
23     x = f(&arg);
24
25     printf("x: %lu\n", x);
26     printf("dereference q: %lu\n", *q);
27
28     return 0;
29 }
30
31 uint64_t f(uint64_t* ptr)
32 {
33     uint64_t x = 0;
34     x = g(*ptr);
35     return x + 1;
36 }
37
38 uint64_t g(uint64_t a)
39 {
40     uint64_t x = 2*a;
41     q = &x; // <-- THIS IS AN ERROR (AKA BUG)
42     return x;
43 }
```



-----[as.txt (x86)]-----

```

1  2. A look at the assembly...
2
3  To see the assembly code that the C compiler (gcc) produces:
4  $ gcc -O0 -S example.c
5  (then look at example.s.)
6  NOTE: what we show below is not exactly what gcc produces. We have
7  simplified, omitted, and modified certain things.
8
9  main:
10    pushq %rbp          # prologue: store caller's frame pointer
11    movq %rsp, %rbp      # prologue: set frame pointer for new frame
12
13    → subq $16, %rsp    # make stack space
14
15    → movq $0, -8(%rbp) # x = 0 (x lives at address rbp - 8)
16    → movq $8, -16(%rbp) # arg = 8 (arg lives at address rbp - 16)
17
18    → leaq -16(%rbp), %rdi # load the address of (rbp-16) into %rdi
19    →           # this implements "get ready to pass (&arg) to f"
20
21    → call f            # invoke f
22
23    → movq %rax, -8(%rbp) # x = (return value of f)
24
25    # eliding the rest of main()
26
27 f:
28    pushq %rbp          # prologue: store caller's frame pointer
29    movq %rsp, %rbp      # prologue: set frame pointer for new frame
30
31    → subq $32, %rsp    # make stack space
32    → movq %rdi, -24(%rbp) # Move ptr to the stack
33    →           # (ptr now lives at rbp - 24)
34    → movq $0, -8(%rbp) # x = 0 (x's address is rbp - 8)
35
36    → movq -24(%rbp), %r8 # move 'ptr' to %r8
37    → movq (%r8), %r9   # dereference 'ptr' and save value to %r9
38    → movq %r9, %rdi    # Move the value of *ptr to rdi,
39    →           # so we can call g
40
41    call g              # invoke g
42
43    → movq %rax, -8(%rbp) # x = (return value of g)
44    → movq -8(%rbp), %r10 # compute x + 1, part I
45    → addq $1, %r10       # compute x + 1, part II
46    → movq %r10, %rax    # Get ready to return x + 1
47
48    → movq %rbp, %rsp    # epilogue: undo stack frame
49    → popq %rbp          # epilogue: restore frame pointer from caller
50    → ret                # return
51
52
53 g:
54    pushq %rbp          # prologue: store caller's frame pointer
55    movq %rsp, %rbp      # prologue: set frame pointer for new frame
56
57    ....
58
59    movq %rbp, %rsp      # epilogue: undo stack frame
60    popq %rbp            # epilogue: restore frame pointer from caller
61    ret
```

Register	Usage	Preserved across function calls
%rax = 8 %rbx = 9 %rcx = 10 %rdx = 11	temporary register; with variable arguments passes information about the number of vector registers used; 1 st return register callee-saved register used to pass 4 th integer argument to functions used to pass 3 rd argument to functions; 2 nd return register	No
%rsp %rbp	stack pointer callee-saved register; optionally used as frame pointer	Yes Yes
%rsi %rdi	used to pass 2 nd argument to functions used to pass 1 st argument to functions	No No
%r8 %r9 %r10	used to pass 5 th argument to functions used to pass 6 th argument to functions temporary register, used for passing a function's static chain pointer	No No No
%r11 %r12-r14 %r15	temporary register callee-saved registers callee-saved register; optionally used as GOT base pointer	No Yes Yes

Borrowed from
<https://github.com/hjl-tools/x86-psABI/wiki/x86-64-psABI-1.0.pdf>
Figure 3.4

