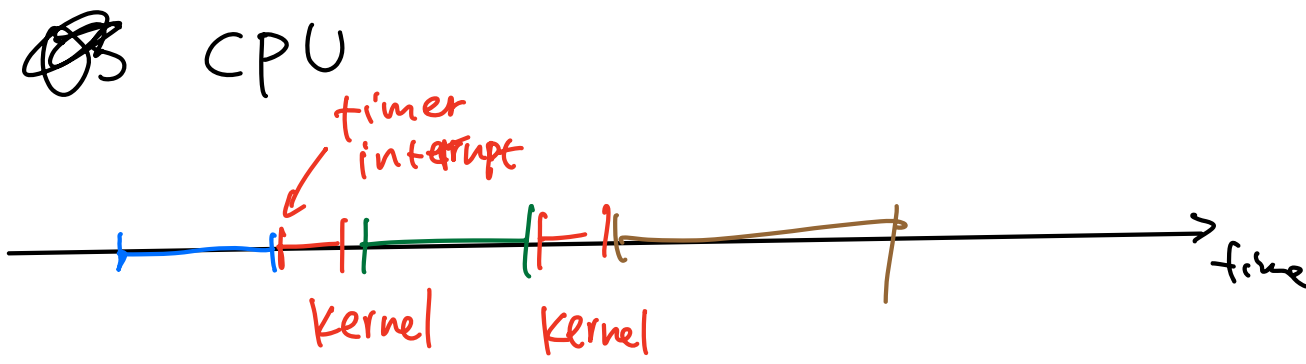


Week 4
CS4973/CS6640
01/27 2025
<https://naizhengtan.github.io/25spring/>

- ✓ 1. Timer interrupt ← (brief)
- ✓ 2. Review: OS scheduling ←
- ✓ 3. egos scheduler & lab3

lab2



Timer interrupts

A diagram with the text 'Timer interrupts' in a large, bold, black font. To the right of this text, there are two hand-drawn arrows. The upper arrow curves upwards and to the right, pointing towards the text 'CPU'. The lower arrow curves downwards and to the right, pointing towards the text 'Programmer'. Both 'CPU' and 'Programmer' are written in a casual, handwritten style.

[borrowed from Yunhao's
CS 4411/5411: Practicum in Operating Systems, 22fall]

CPU view: Core-local Interrupt (CLINT)

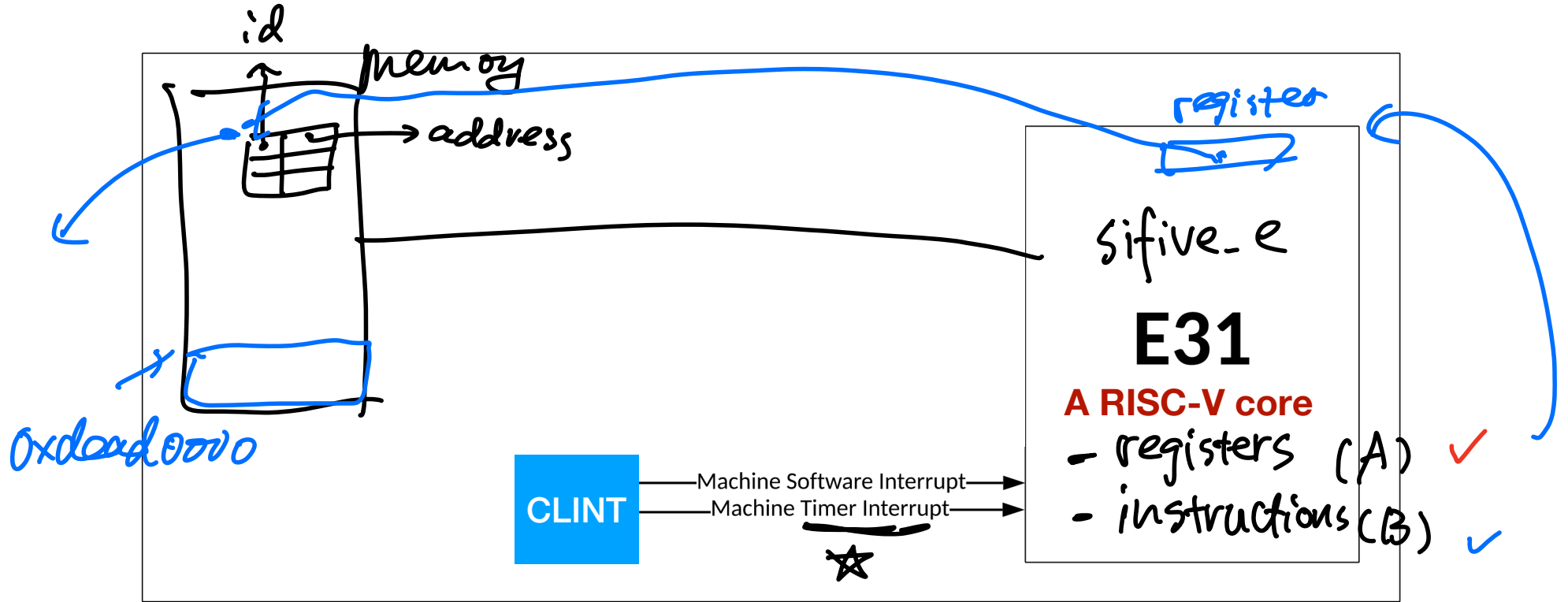


Figure 4 of Sifive FE310 manual

A programmer's view

```
void handler() {  
    printf("Got a timer interrupt!");  
    // (4) reset timer  
}  
  
int main() {  
    // (1) register interrupt handler  
    // (2) set a timer  
    // (3) enable timer interrupt  
  
    while(1);  
}
```

Q: What
if you
were
CPU designer



machine-mode

The mtvec CSR

handler's address →
trap →



Value	Name	Description
0	Direct	All exceptions set pc to BASE.
1	Vectored	Asynchronous interrupts set pc to BASE+4×cause.
≥2	—	Reserved

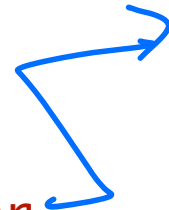
Table 3.5: Encoding of mtvec MODE field.

A programmer's view

```
void handler() {  
    printf("Got a timer interrupt!");  
    // (4) reset timer  
}  
  
int main() {  
    // (1) register interrupt handler  
    // (2) set a timer  
    // (3) enable timer interrupt  
  
    while(1);  
}
```

put address
to %mtvec.

0x810000C



OSI Handout Week4

1. Timer interrupt handler

```

void handler() {
    CRITICAL("Got a timer interrupt!");
    // (4) reset timer

}

int main() {
    CRITICAL("This is a simple timer example");

    // (1) register handler() as interrupt handler

    // (2) set a timer

    // (3) enable timer interrupt

    while(1);
}

```

2. Background: RISC-V assembly II

Assembler instructions with C expression operands:

`asm(Template : OutputOperands : InputOperands)`

a) **Template**: a string that is the template for the assembler code.

→ `asm("mret");` **ret**

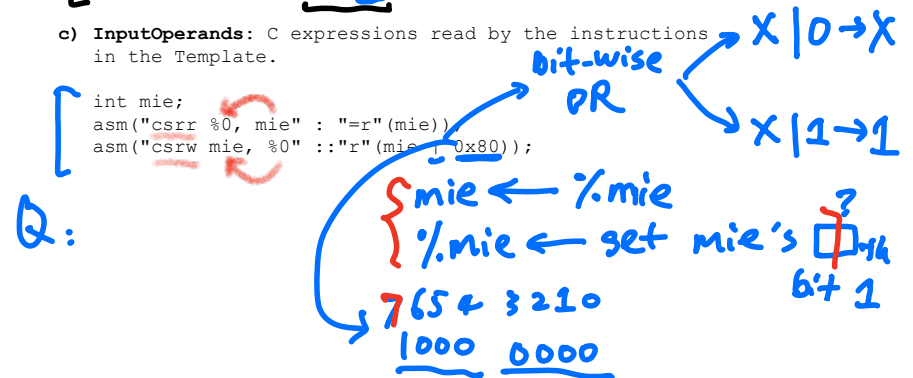
b) **OutputOperands**: the C variables modified by the instructions in the Template.

`void *sp;`
`asm("mv %0, sp" : "=r" sp);`

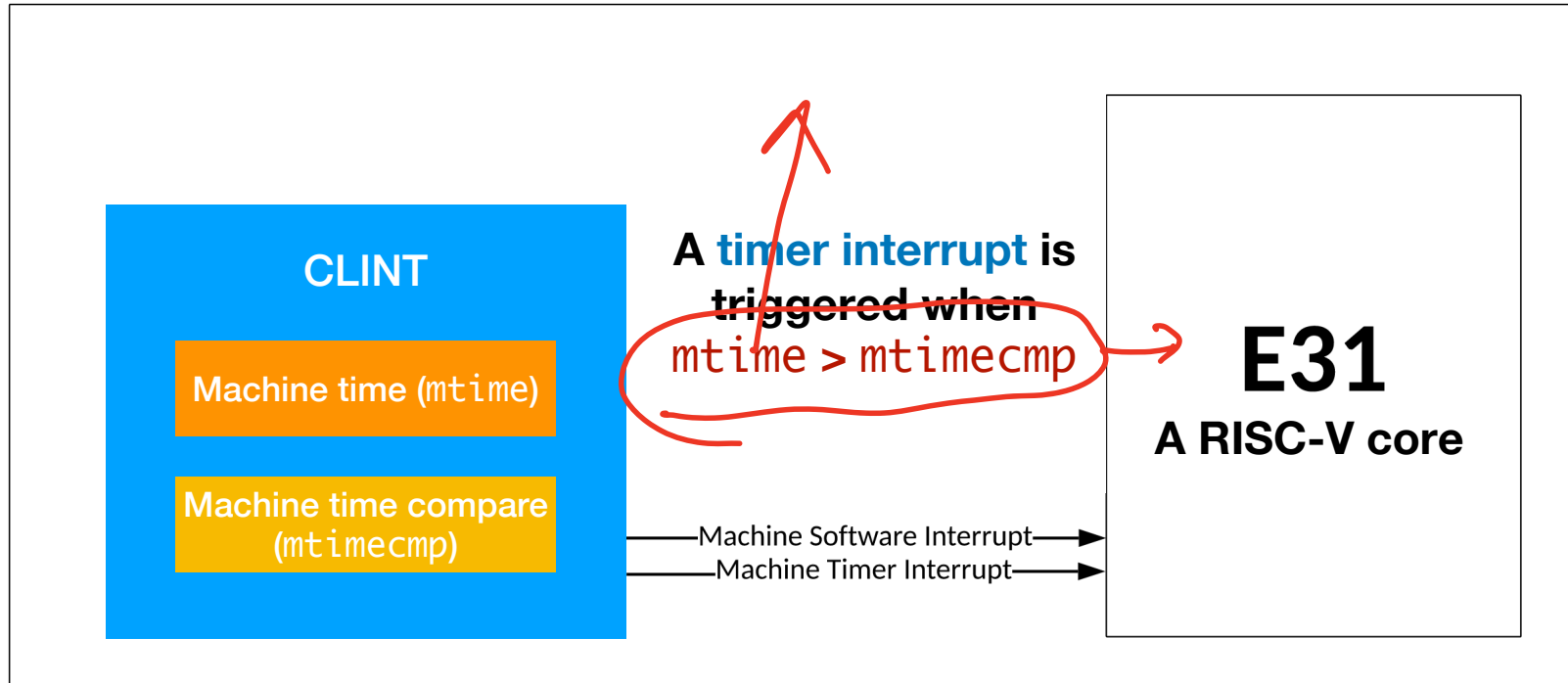
c) **InputOperands**: C expressions read by the instructions in the Template.

`int mie;`
`asm("csrr %0, mie" : "=r"(mie));`
`asm("csrw mie, %0" : "r"(mie & 0x80));`

Q:





The `mtime` and `mtimecmp` CSRs

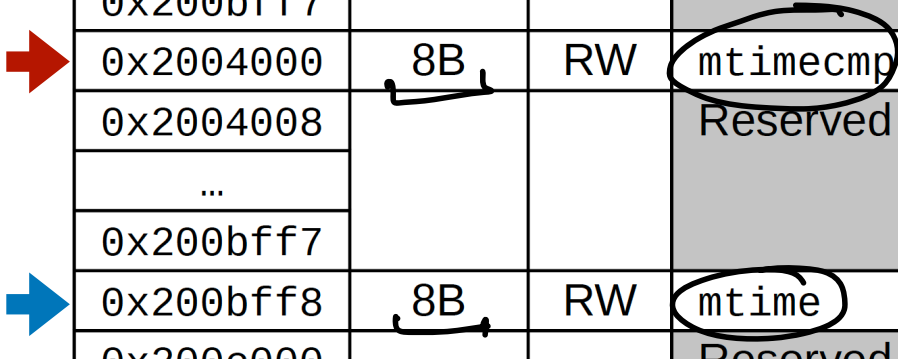


CLINT CSRs are **memory-mapped**

Address	Width	Attr.	Description
0x20000000	4B	RW	msip for hart 0
0x2004008			Reserved
...			
0x200bff7			
0x2004000	8B	RW	mtimecmp for hart 0
0x2004008			Reserved
...			
0x200bff7			
0x200bff8	8B	RW	mtime
0x200c000			Reserved

`mtimecmp_set()` writes 8 bytes to 

`mtime_get()` reads 8 bytes from 



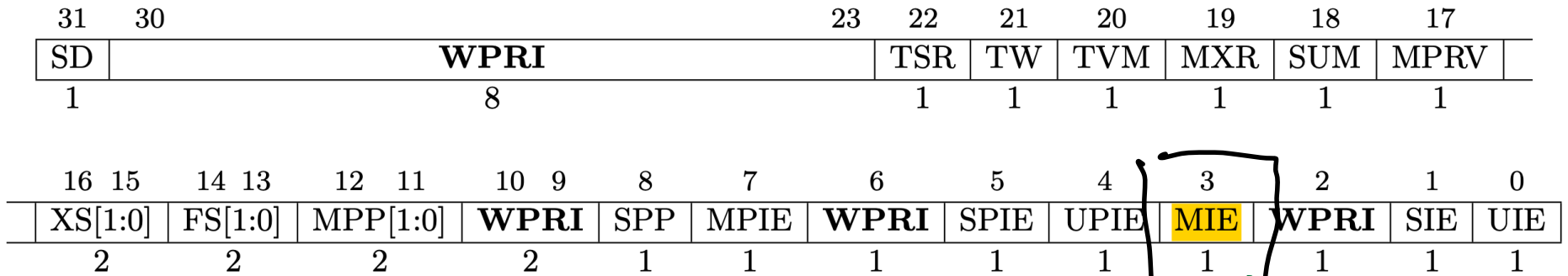
A programmer's view

```
void handler() {  
    printf("Got a timer interrupt!");  
    // (4) reset timer  
}
```

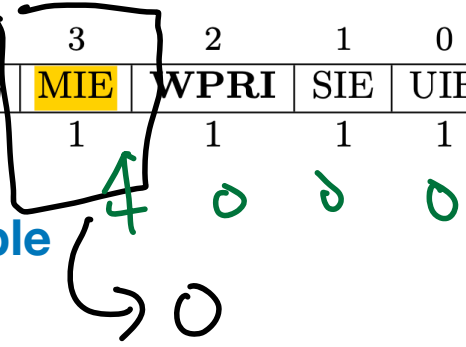
```
int main() {  
    // (1) register interrupt handler ✓  
    // (2) set a timer ✓  
    // (3) enable timer interrupt  
  
    while(1);  
}
```

← ? FZLLN

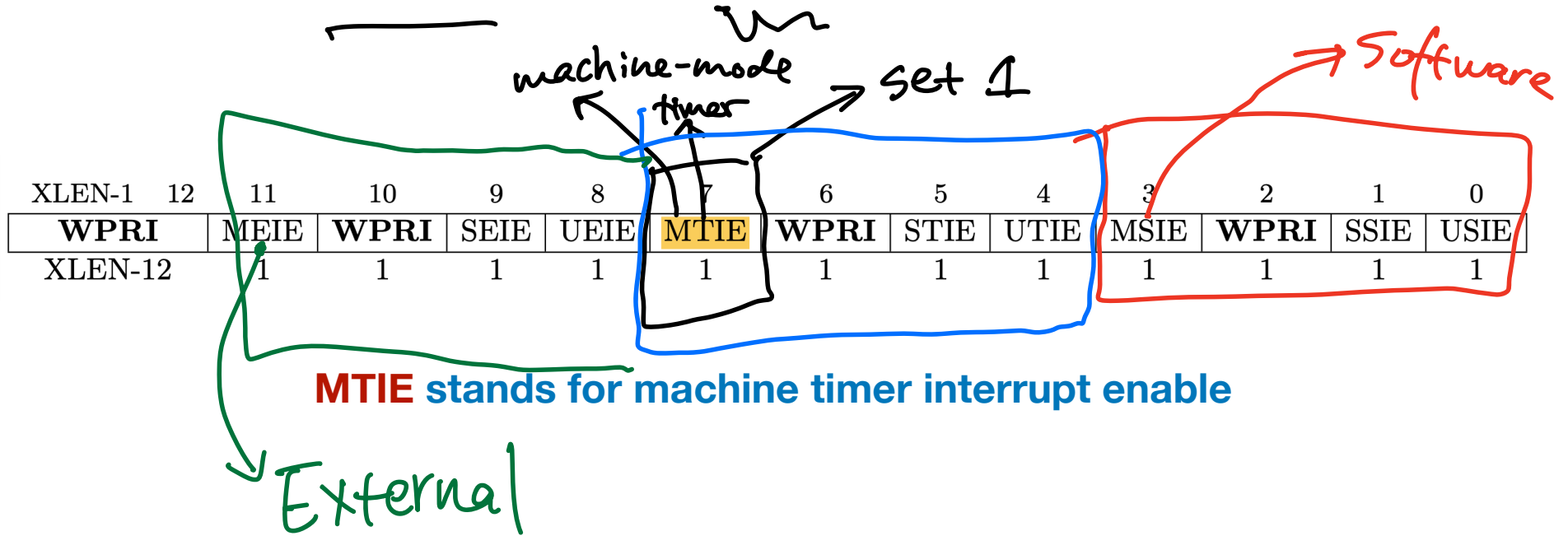
The **mstatus** CSR



MIE stands for machine interrupt enable



The mie CSR (not mstatus.MIE)



• Review

RR → ls
 loop &

MLFQ

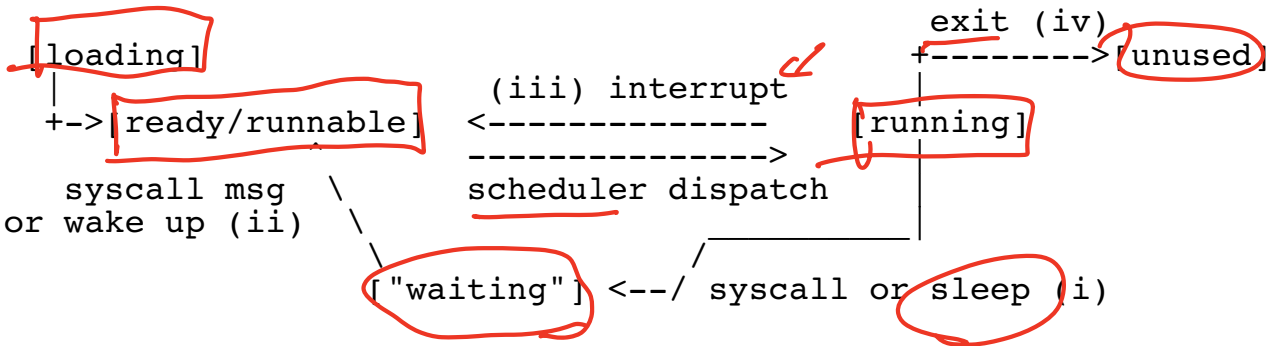
new process



high
mid
low

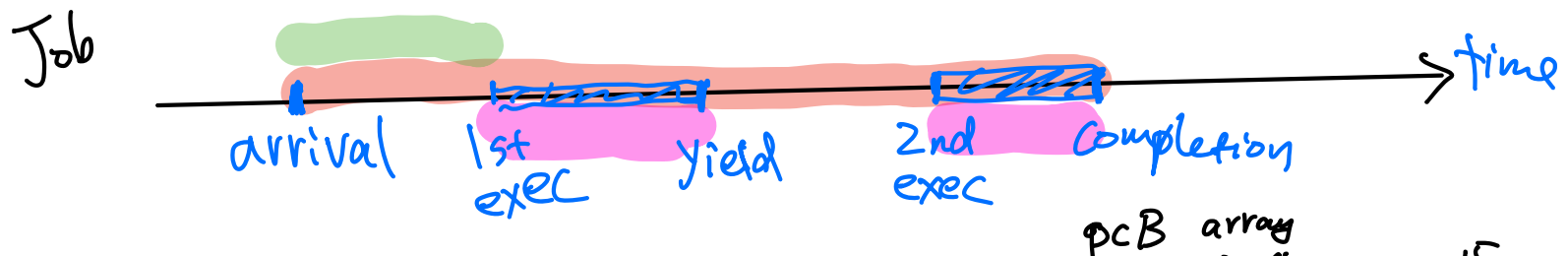
① WHY?

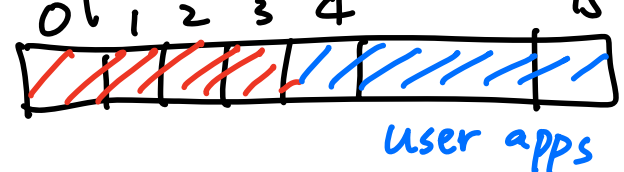
② HOW?



• metrics

- turnaround time : (completion - arrival)
- response time : (1st exec - arrival)
- CPU time : CPU consumed time





3. Machine-mode exception CSRs

- a) **mstatus**
Machine Status, holds the global interrupt enable, along with a plethora of other state.
- b) **mie**
Machine Interrupt Enable, lists which interrupts the processor can take and which it must ignore
- c) **mcause**
Machine Exception Cause, indicates which exception occurred
- d) **mtvec**
Machine Trap Vector, holds the address the processor jumps to when an exception occurs
- e) **mepc**
Machine Exception PC, points to the instruction where the exception occurred
- f) **mtval**
Machine Trap Value, holds additional trap information: the faulting address for address exceptions, the instruction itself for illegal instruction exceptions, and zero for other exceptions
- g) **mip**
Machine Interrupt Pending, lists the interrupts currently pending

4. egos process management (sifive_e)

a) process control block (PCB)

```
[grass/process.h]
struct process {
  int pid;
  int status;
  int receiver_pid; /* used when waiting to send a message */
  void *sp, *mepc; /* process context = stack pointer (sp)
                  * + machine exception program counter (mepc) */
  // scheduling attributes
  union {
    unsigned char chars[64];
    unsigned int ints[16];
    float floats[16];
    unsigned long long longlongs[8];
    double doubles[8];
  } sched_attr;
};
```

Handwritten notes: pid ↔ index in proc-set [16]. Arrows point from 'pid' and 'index' to the 'pid' field in the struct. A vertical list '1 2 3 4 5 6' is next to 'index', and '0 1' is next to 'pid'. '64B' is written next to the union block.

b) global process data structures

```
[grass/kernel.c]
int proc_curr_idx;
struct process proc_set[MAX_NPROCESS];

[grass/process.h]
#define curr_pid proc_set[proc_curr_idx].pid
#define curr_status proc_set[proc_curr_idx].status
```

Handwritten notes: 'index' with an arrow pointing to 'proc_set'. 'BH=0' with an arrow pointing to the 'proc_set' array. A blue bracket underlines the #define lines.

c) process life cycles

```
[grass/scheduler.c]
life-cycle functions:
* proc_on_arrive(int pid): when creating pid
* proc_yield(): when deciding next running process
* proc_on_stop(int pid): when destroying pid
```

a process's life cycle:

- proc_on_arrive() ->
- proc_yield() -> [other proc] -> [ctx_switch to this proc] ->
- proc_yield() -> [other proc] -> [ctx_switch to this proc] ->
- ...
- -> proc_on_stop()

Handwritten diagrams: A tree diagram shows a root box containing '0 0 0 X' with arrows pointing to four child boxes. Below the children are '0x41 0 0 0' and 'A '10''. A red arrow points from the root to '0 0x41 0 0'. Below that is 'NOTHING' in red. At the bottom, a green box contains '0x42 0x41 0 0' and 'B A '10''.

Next...

- two bummers in setting timers
- interrupts beyond CLINT
- RISC-V fine-grained control of interrupts
- how to know which interrupt is triggered?

a bummer: mtime rollover

Higher 4 bytes

Lower 4 bytes

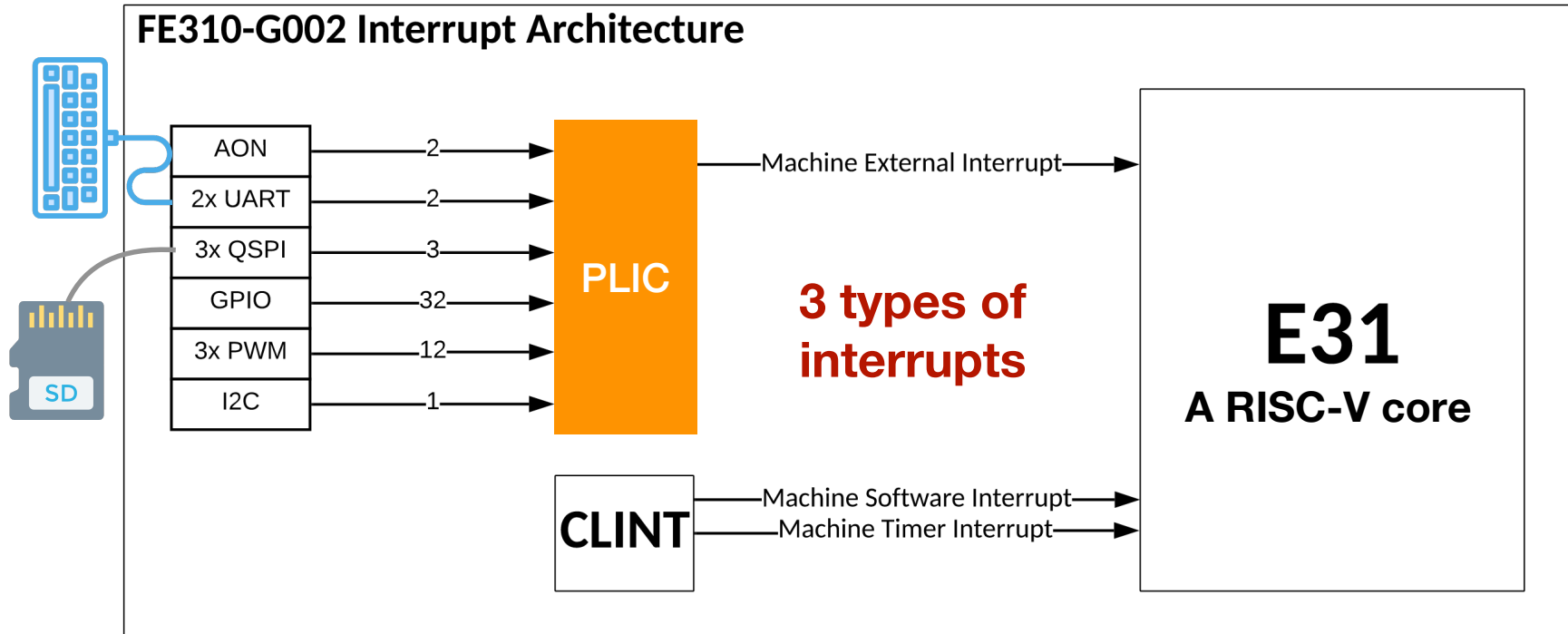
When reading the **lower** 4 bytes, mtime is
`0x00000000` `0xffffffff`

When reading the **higher** 4 bytes, mtime is
`0x00000001` `0x00000000`

Combine the two `0x00000001ffffffff` is wrong!

There is another bummer: mtime overflow

Platform Interrupt Controller (PLIC)



mie provides fine-grained control

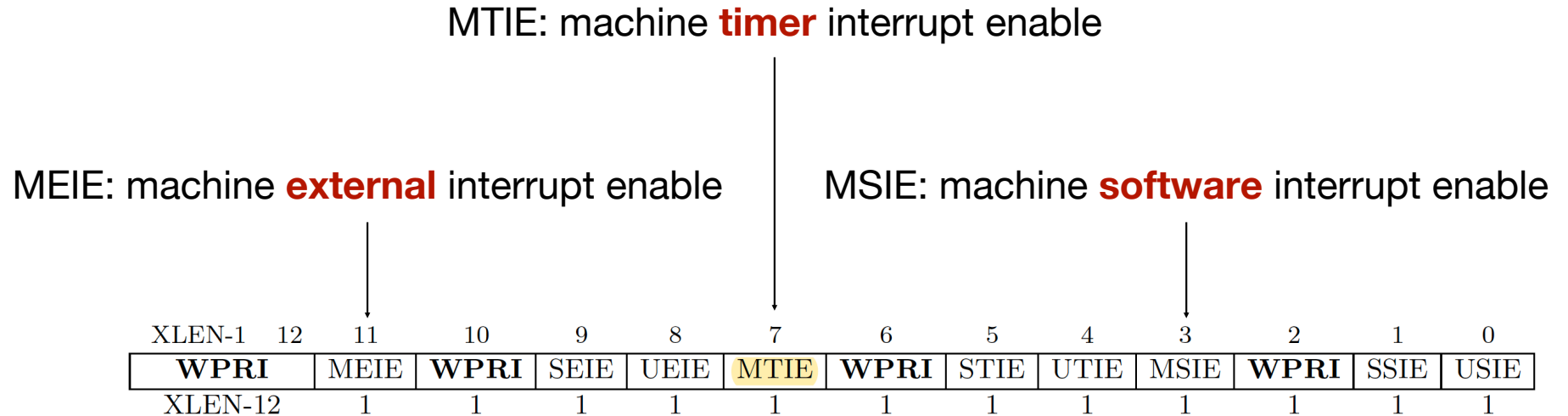


Figure 3.12: Machine interrupt-enable register (`mie`).

Timer is interrupt #7

Interrupts

Exceptions

Interrupt	Exception Code	Description
1	0	<i>Reserved</i>
1	1	Supervisor software interrupt
1	2	<i>Reserved</i>
1	3	Machine software interrupt
1	4	<i>Reserved</i>
1	5	Supervisor timer interrupt
1	6	<i>Reserved</i>
1	7	Machine timer interrupt
1	8	<i>Reserved</i>
1	9	Supervisor external interrupt
1	10	<i>Reserved</i>
1	11	Machine external interrupt
1	12-15	<i>Reserved</i>
1	≥ 16	<i>Designated for platform use</i>
0	0	Instruction address misaligned
0	1	Instruction access fault
0	2	Illegal instruction
0	3	Breakpoint
0	4	Load address misaligned
0	5	Load access fault
0	6	Store/AMO address misaligned
0	7	Store/AMO access fault
0	8	Environment call from U-mode
0	9	Environment call from S-mode
0	10	<i>Reserved</i>
0	11	Environment call from M-mode
0	12	Instruction page fault
0	13	Load page fault
0	14	<i>Reserved</i>
0	15	Store/AMO page fault
0	16-23	<i>Reserved</i>
0	24-31	<i>Designated for custom use</i>
0	32-47	<i>Reserved</i>
0	48-63	<i>Designated for custom use</i>
0	≥ 64	<i>Reserved</i>



System calls?

Interrupts

Exceptions

Interrupt	Exception Code	Description
1	0	<i>Reserved</i>
1	1	Supervisor software interrupt
1	2	<i>Reserved</i>
1	3	Machine software interrupt
1	4	<i>Reserved</i>
1	5	Supervisor timer interrupt
1	6	<i>Reserved</i>
1	7	Machine timer interrupt
1	8	<i>Reserved</i>
1	9	Supervisor external interrupt
1	10	<i>Reserved</i>
1	11	Machine external interrupt
1	12-15	<i>Reserved</i>
1	≥ 16	<i>Designated for platform use</i>
0	0	Instruction address misaligned
0	1	Instruction access fault
0	2	Illegal instruction
0	3	Breakpoint
0	4	Load address misaligned
0	5	Load access fault
0	6	Store/AMO address misaligned
0	7	Store/AMO access fault
0	8	Environment call from U-mode
0	9	Environment call from S-mode
0	10	<i>Reserved</i>
0	11	Environment call from M-mode
0	12	Instruction page fault
0	13	Load page fault
0	14	<i>Reserved</i>
0	15	Store/AMO page fault
0	16-23	<i>Reserved</i>
0	24-31	<i>Designated for custom use</i>
0	32-47	<i>Reserved</i>
0	48-63	<i>Designated for custom use</i>
0	≥ 64	<i>Reserved</i>

